

Nexus 7000配置和验证LISP IGP辅助扩展子网模式

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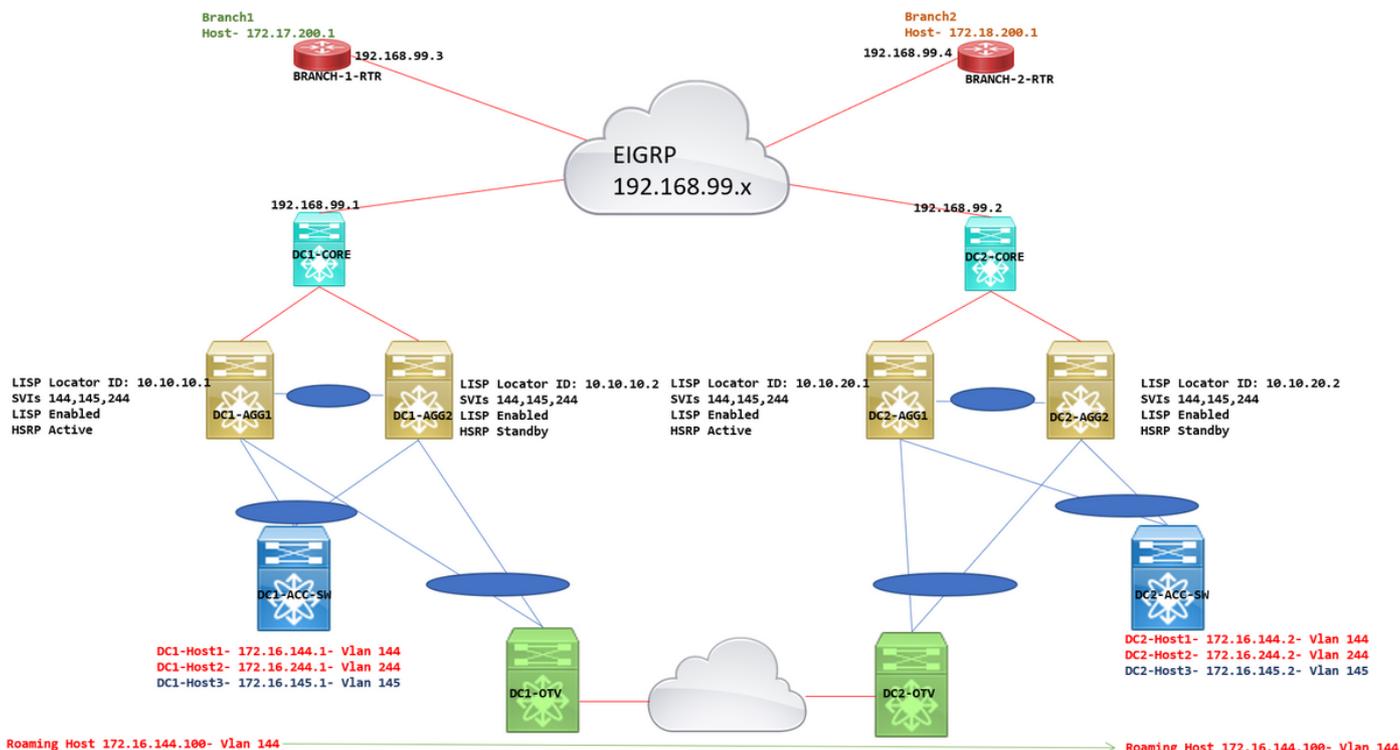
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简介

本文档将介绍如何使用Nexus 7000部署LISP IGP辅助扩展子网模式(ESM)

拓扑



拓扑详细信息

- DC1和DC2是OTV扩展的两个位置
- VLAN 144、145和244在所有Agg、接入层和OTV交换机上配置
- 这些VLAN的SVI在Agg交换机上配置。SVI 144和244位于VRF租户-1中；SVI 145在VRF租户2中。
- 在部署LISP IGP Assist时，SVI不必位于VRF中；本示例使用多个VRF只是为了说明所需的配置更改（在每个相关VRF环境下）；所有SVI都可以位于同一VRF中，并且仍然可以使用LISP IGP辅助
- HSRP在Vlan144、145和244中配置；此拓扑中配置了FHRP隔离，这意味着总共4台交换机将运行HSRP，并且两端都有主用/备用对。FHRP隔离通过过滤HSRP Hello消息来实现。
- DC1-agg1和DC2-Agg2是vPC对；DC2-Agg1和DC2-Agg2也适用
- LISP配置在SVI 144、145和244下应用
- 每个VRF从Agg建立到核心交换机的EIGRP邻居关系。子接口从每个VRF的Agg交换机运行到核心交换机，并且EIGRP邻居关系形成在这些子接口上。
- 远程路由器（分支）也属于同一IGP域。
- 使用LISP IGP Assist时，没有LISP封装/解封，因此LISP路由必须重分发到IGP（此处为EIGRP）。对于本文档中描述的此部署模式，分支路由器将没有任何LISP配置。

使用的组件

- Nexus 7000是核心交换机，带SUP2E，F3/M3运行8.2(4)NXOS版本
- 分支路由器是ASR1k
- OTV在这些Nexus 7000交换机的另一个VDC中配置；OTV和LISP必须位于不同的VDC上。共享VDC不是选项。

本文档中的信息都是基于特定实验室环境中的设备编写的。本文档中使用的所有设备最初均采用原始（默认）配置。如果您使用的是真实网络，请确保您已经了解所有命令的潜在影响。

AGG交换机上所需的配置

DC1-Agg1和DC1-Agg2上的LISP特定配置

Common Configuration on both DC1-Agg1 and DC1-Agg2

```
feature lisp
vrf context tenant-1                                     # This example is
based on SVI 144 in VRF- tenant-1 and SVI 145 in VRF- tenant-2
  ip lisp etr                                           # This is needed to
initialize LISP and only etr is needed on a IGP assist mode Environment
  lisp instance-id 2                                    # Instance-ID should
be unique per VRF
  ip lisp locator-vrf default                            # Locator Is
specified in Default VRF
  lisp dynamic-eid VLAN144                              # Dynamic EID
definition for Vlan 144
  database-mapping 172.16.144.0/24 10.10.10.1 priority 50 weight 50 # Database-mapping
for 172.16.144.0/24 which is the Vlan 144; IP-> 10.10.10.1 is the Loopback100 IP address(which
is the unique IP on DC1-AGG1)
  database-mapping 172.16.144.0/24 10.10.10.2 priority 50 weight 50 # Database-mapping
for 172.16.144.0/24 which is the Vlan 144; IP-> 10.10.10.2 is the Loopback100 IP address(which
is the unique IP on DC1-AGG2)
  map-notify-group 239.254.254.254                      # Multicast group
that will be used by LISP enabled switches to communicate about new EID learns or periodic EID
notification messages
  no route-export away-dyn-eid                          # This is a hidden
command required to stop advertising any null0 /32 route for a remote host to the IGP
  lisp dynamic-eid VLAN244                              # Dynamic EID
definition for Vlan 244
  database-mapping 172.16.244.0/24 10.10.10.1 priority 50 weight 50
  database-mapping 172.16.244.0/24 10.10.10.2 priority 50 weight 50
  map-notify-group 239.254.254.254
  no route-export away-dyn-eid

vrf context tenant-2
  ip lisp etr
  lisp instance-id 3
  ip lisp locator-vrf default
  lisp dynamic-eid VLAN145
  database-mapping 172.16.145.0/24 10.10.10.1 priority 50 weight 50
  database-mapping 172.16.145.0/24 10.10.10.2 priority 50 weight 50
  map-notify-group 239.254.254.254
  no route-export away-dyn-eid
```

Configuration on DC1-Agg1

```
interface Vlan144
  no shutdown
  vrf member tenant-1
  lisp mobility VLAN144
  lisp extended-subnet-mode                             # SVI needs to be in
ESM Mode-Extended subnet mode
  ip address 172.16.144.250/24
  ip pim sparse-mode
  hsrp 144
  preempt
  priority 254
  ip 172.16.144.254
```

```
interface Vlan145
  no shutdown
  vrf member tenant-2
  lisp mobility VLAN145
  lisp extended-subnet-mode
  ip address 172.16.145.250/24
  ip pim sparse-mode
  hsrp 145
    preempt
    priority 254
    ip 172.16.145.254
```

```
interface Vlan244
  no shutdown
  vrf member tenant-1
  lisp mobility VLAN244
  lisp extended-subnet-mode
  ip address 172.16.244.250/24
  hsrp 244
    preempt
    priority 254
    ip 172.16.244.254
```

```
interface loopback100
  ip address 10.10.10.1/32
  ip router eigrp 100
  ip pim sparse-mode
```

Configuration on DC1-Agg2

```
interface Vlan144
  no shutdown
  vrf member tenant-1
  lisp mobility VLAN144
  lisp extended-subnet-mode
  ip address 172.16.144.251/24
  ip pim sparse-mode
  hsrp 144
    ip 172.16.144.254
```

```
interface Vlan145
  no shutdown
  vrf member tenant-2
  lisp mobility VLAN145
  lisp extended-subnet-mode
  ip address 172.16.145.251/24
  ip pim sparse-mode
  hsrp 145
    ip 172.16.145.254
```

```
interface Vlan244
  no shutdown
  vrf member tenant-1
  lisp mobility VLAN244
  lisp extended-subnet-mode
  no ip redirects
  ip address 172.16.244.251/24
  hsrp 244
    ip 172.16.244.254
```

```
interface loopback100
  ip address 10.10.10.2/32
  ip router eigrp 100
  ip pim sparse-mode
```

#数据库映射的提供方式必须是，在一端，DC1-Agg1和DC1-Agg2环回IP地址都需要指定；在DC2-Agg1和DC2-Agg2中，必须创建一个唯一环回，并将其放在数据库映射中。

#在IGP辅助模式下，如果使用配置 —> "ip lisp itr-etr"，则会为未启用LISP的VLAN注入/32 null0主机路由；因此，IGP辅助模式的正确配置是“ip lisp etr”。

DC2-Agg1和DC2-Agg2上的LISP特定配置

Common Configuration on both DC2-Agg1 and DC2-Agg2

```
feature lisp

vrf context tenant-1
  ip lisp etr
  lisp instance-id 2
  ip lisp locator-vrf default
  lisp dynamic-eid VLAN144
    database-mapping 172.16.144.0/24 10.10.20.1 priority 50 weight 50      # Note that the IP
addresses used in DC2 Agg switches are 10.10.20.1 and 10.10.20.2(Which are Loopbacks Configured
on DC2-Agg switches)
    database-mapping 172.16.144.0/24 10.10.20.2 priority 50 weight 50
    map-notify-group 239.254.254.254
    no route-export away-dyn-eid
  lisp dynamic-eid VLAN244
    database-mapping 172.16.244.0/24 10.10.20.1 priority 50 weight 50
    database-mapping 172.16.244.0/24 10.10.20.2 priority 50 weight 50
    map-notify-group 239.254.254.254
    no route-export away-dyn-eid
vrf context tenant-2
  ip lisp etr
  lisp instance-id 3
  ip lisp locator-vrf default
  lisp dynamic-eid VLAN145
    database-mapping 172.16.145.0/24 10.10.20.1 priority 50 weight 50
    database-mapping 172.16.145.0/24 10.10.20.2 priority 50 weight 50
    map-notify-group 239.254.254.254
    no route-export away-dyn-eid
```

Configuration on DC2-Agg1

```
interface Vlan144 no shutdown vrf member tenant-1 lisp mobility VLAN144 lisp extended-subnet-
mode ip address 172.16.144.252/24 ip pim sparse-mode hsrp 144 preempt priority 254 ip
172.16.144.254 interface Vlan145 no shutdown vrf member tenant-2 lisp mobility VLAN145 lisp
extended-subnet-mode ip address 172.16.145.252/24 ip pim sparse-mode hsrp 145 preempt priority
254 ip 172.16.145.254 interface Vlan244 no shutdown vrf member tenant-1 lisp mobility VLAN244
lisp extended-subnet-mode ip redirects ip address 172.16.244.252/24 hsrp 244 preempt priority
254 ip 172.16.244.254 interface loopback100 ip address 10.10.20.1/32 ip router eigrp 100 ip pim
sparse-mode
```

```
Configuration on DC2-Agg2
interface Vlan144 no shutdown vrf member tenant-1 lisp mobility VLAN144 lisp extended-subnet-
mode ip address 172.16.144.253/24 ip pim sparse-mode hsrp 144 ip 172.16.144.254 interface
Vlan145 no shutdown vrf member tenant-2 lisp mobility VLAN145 lisp extended-subnet-mode ip
address 172.16.145.253/24 ip pim sparse-mode hsrp 145 ip 172.16.145.254 interface Vlan244 no
shutdown vrf member tenant-1 lisp mobility VLAN244 lisp extended-subnet-mode no ip redirects ip
address 172.16.244.253/24 hsrp 244 preempt ip 172.16.244.254 interface loopback100 ip address
10.10.20.2/32 ip router eigrp 100 ip pim sparse-mode
```

DC1和DC2 Agg LISP配置之间的区别是“数据库映射”中定义的环回。在DC1配置中，这将使用DC1-Agg1和DC1-Agg2的环回定义，对于DC2，数据库映射将使用DC2-Agg1和DC2-Agg2中的环回定义

#下面显示的其余IGP/路由映射/前缀列表配置将相似（为接口分配的IP地址确实不同）

IGP特定

```
router eigrp 100
  address-family ipv4 unicast
  vrf tenant-1
    distance 90 245 # External EIGRP
Routes have to have an AD which is higher than the default LISP AD(which is 240); Reason being,
if the redistributed route from dc1-aggr1 comes back to dc1-aggr2 via eigrp, default EIGRP
External is 170 which will override LISP route causing problems
    redistribute lisp route-map lisp-to-eigrp # This command is to
  redistribute LISP /32 routes only to the IGP(EIGRP In this example)
    redistribute direct route-map direct # This is needed so
that the direct routes(/24 SVI routes in LISP) are redistributed to the IGP; This will be needed
if there is some device that is trying to communicate to a silent host in the LISP enabled Vlan
  vrf tenant-2
    distance 90 245
    redistribute lisp route-map lisp-to-eigrp
    redistribute direct route-map direct
```

#启用LISP的AGG VDC还将与核心端形成IGP邻居关系

#在本示例中，作为每个租户VRF一部分的子接口用于形成与核心的邻居关系，如下所示。

```
interface Ethernet3/6.111
  encapsulation dot1q 111
  vrf member tenant-1
  ip address 192.168.98.1/30
  ip router eigrp 100
  no shutdown
```

```
interface Ethernet3/6.212
  encapsulation dot1q 212
  vrf member tenant-2
  ip address 192.168.198.1/30
  ip router eigrp 100
  no shutdown
```

路由映射/前缀列表

```
ip prefix-list lisp-to-eigrp seq 5 permit 0.0.0.0/0 ge 32 # This is the prefix
list that is matching any /32 routes which are to be redistributed from LISP To IGP

route-map direct permit 10 # This is for the
Direct routes

route-map lisp-to-eigrp deny 10 # This is to prevent
any null0 routes from being redistributed to IGP from LISP
  match interface Null0

route-map lisp-to-eigrp permit 20 # This is to allow
redistribution of /32 host routes
  match ip address prefix-list lisp-to-eigrp
```

#所有AGG交换机 (DC1和DC2) 都需要上述所有配置。请记住，为SVI、环回、HSRP VIP提供唯一的IP地址对于所有SVI都是相同的

OTV VDC配置

HSRP过滤

#对于IGP辅助部署，当通过OTV或任何其他机制进行扩展时，必须进行FHRP隔离；

#这通过过滤OTV VDC中的FHRP Hello消息来完成

#在本示例中，使用N7k OTV，因此应用了以下配置来过滤OTV VDC中的FHRP数据包。

```
ip access-list ALL_IPs
 10 permit ip any any
mac access-list ALL_MACs
 10 permit any any
ip access-list HSRP_IP
 10 permit udp any 224.0.0.2/32 eq 1985
 20 permit udp any 224.0.0.102/32 eq 1985
mac access-list HSRP_VMAC
 10 permit 0000.0c07.ac00 0000.0000.00ff any
 20 permit 0000.0c9f.f000 0000.0000.0fff any
arp access-list HSRP_VMAC_ARP
 10 deny ip any mac 0000.0c07.ac00 ffff.ffff.ff00
 20 deny ip any mac 0000.0c9f.f000 ffff.ffff.f000
 30 permit ip any mac any
vlan access-map HSRP_Localization 10
  match mac address HSRP_VMAC
  match ip address HSRP_IP
  action drop
vlan access-map HSRP_Localization 20
  match mac address ALL_MACs
  match ip address ALL_IPs
  action forward
vlan filter HSRP_Localization vlan-list 144-145
ip arp inspection filter HSRP_VMAC_ARP vlan 144-145

mac-list OTV_HSRP_VMAC_deny seq 10 deny 0000.0c07.ac00 ffff.ffff.ff00
mac-list OTV_HSRP_VMAC_deny seq 11 deny 0000.0c9f.f000 ffff.ffff.f000
mac-list OTV_HSRP_VMAC_deny seq 20 permit 0000.0000.0000 0000.0000.0000
route-map OTV_HSRP_filter permit 10
  match mac-list OTV_HSRP_VMAC_deny

otv-isis default
  vpn Overlay0
  redistribute filter route-map OTV_HSRP_filter
```

#仅OTV VDC上需要FHRP过滤配置；如果使用ASR OTV部署，则过滤机制应作为相关机制使用，并根据ASR配置指南进行记录。

OTV抑制ARP

#在OTV VDC上禁用ARP ND缓存功能

```
interface Overlay0
  no otv suppress-arp-nd >>>>>
```

由于LISP配置导致的路由填充

```

DC1-AGG1# show ip route lisp vrf tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
'***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>

172.16.144.0/25, ubest/mbest: 1/0
    *via Null0, [240/1], 07:22:30, lisp, dyn-eid
172.16.144.128/25, ubest/mbest: 1/0
    *via Null0, [240/1], 07:22:30, lisp, dyn-eid

```

#在SVI 144上启用LISP后，将自动创建两条Null0路由；SVI 144是/24子网，因此第1条null0路由来自172.16.144.0/25，第二条路由来自172.16.144.128/25，如上所示。

#这是预期的，而且是按设计的；这样做是为了确保源自未发现主机的数据包触发RPF异常，该异常将导致数据包被传送到CPU，并最终有助于主机检测(EID)

当主机在启用LISP的SVI内联机时的事件顺序

#启用LISP的接口上的主机检测基于从数据库映射配置中指定的范围内的IP地址接收的L3流量。

为便于检测主机，请注意，当在接口上启用LISP时：

- #接口上启用了RPF异常，因此未知源生成的数据包会触发异常
- #源于Null0的LISP路由已安装，以确保未知源触发RPF异常

由于此解决方案依赖OTV在两个数据中心之间进行L2扩展，因此ARP信令不能直接用于检测IP主机，因为在许多情况下，ARP信令会广播到所有交换机。

但是，ARP信号用作LISP指示可能存在未检测到的主机。由于主机可驻留在OTV网桥的任一端，因此LISP在学习新的IP-MAC绑定后启动本地化机制。

定位机制的工作原理如下：

- #交换机获知新的IP-MAC绑定（通过GARP、RARP或ARP请求）。
- #作为活动HSRP工作的交换机向主机发送回应请求，但源自HSRP VIP地址
- #主机响应回应请求，但在OTV中进行FHRP隔离后，仅在主机所在的DC站点上收到回应应答
- #由于回应应答是L3数据包，因此LISP会检测到主机。

#如果在任何启用LISP的SVI上收到IP数据包，则它本身将向LISP进程提供通知终端为本地的信息；不会发送任何ICMP ECHO请求，以进一步确认主机是否是本地主机。因此，务必注意，从DC2主机对DC1-AGG SVI IP地址执行Ping操作会导致端点标识损坏，这也可能导致ping丢失或流量黑洞，因为主机现在被标识为DC1中的本地EID，而不是DC2。因此，ping不应从LISP中的SVI IP地址发起环境，因为这可能损坏路由表，并导致流量黑洞。如果LISP已启用VLAN的主机尝试ping SVI IP地址，则会出现同样的问题；对VIP执行ping操作应该正常，因为两端存在相同且处于活动状态，并且本地站点将捕获数据包。

主机在DC1上线时的路由表条目示例如下：

```

DC1-AGG1# show ip route 172.16.144.1 vrf tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
'***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>

```

```
172.16.144.1/32, ubest/mbest: 1/0, attached
  *via 172.16.144.1, Vlan144, [240/1], 3d05h, lisp, dyn-eid
  via 172.16.144.1, Vlan144, [250/0], 3d05h, am
```

```
DC1-AGG2# sh ip route 172.16.144.1 vr tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
'***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>
```

```
172.16.144.1/32, ubest/mbest: 1/0, attached
  *via 172.16.144.1, Vlan144, [240/1], 3d05h, lisp, dyn-eid
  via 172.16.144.1, Vlan144, [250/0], 3d05h, am
```

#如上所示，有两条路由；AD为250的AM->邻接管理器（由ARP进程填充）使用管理距离为240的LISP进程。

DC1中的两台Agg交换机将具有相同的条目。

#此外，LISP将在动态EID表中列出主机的相同条目，如下所示。

```
DC1-AGG1# show lisp dynamic-eid detail vrf tenant-1 | in 144.1, nex 1 172.16.144.1, Vlan144,
uptime: 3d05h, last activity: 00:14:38 Discovered by: packet reception DC1-AGG2# show lisp
dynamic-eid detail vrf tenant-1 | in 144.1, nex 1 172.16.144.1, Vlan144, uptime: 3d05h, last
activity: 00:00:37 Discovered by: site-based Map-Notify
```

#发现在这两种情况下都不同；HSRP活动的DC1-AGG1通过“数据包接收”记录条目，这基本上意味着传入的数据包导致将其添加为EID

Agg1得知EID后，会从源IP-> Loopback100 IP地址（在数据库映射下定义）向组 — > 239.254.254.254（在上面配置）发送组播消息，vPC对等交换机也会收到该消息并相应地填充条目。由于数据库映射同时具有dc1-agg1和dc1-agg2的IP地址，因此将其视为本地EID。此相同的组播数据包也会通过OTV传输到远程站点；但是，远程站点将检查数据库映射，由于此数据包来自与“数据库映射”不同的IP地址，因此DC2 AGg交换机不会将其视为本地EID。

映射通知消息

#当启用LISP的SVI检测到主机时，触发的“映射通知”消息将发送到在相应动态EID配置下定义的组播组

#除触发的映射通知消息外，该VLAN中有由HSRP活动（或FHRP活动）交换机发送的定期映射通知消息；

#映射通知消息的PCAP如下所示。

```

> Frame 285: 122 bytes on wire (976 bits), 122 bytes captured (976 bits) on interface eth0, id 0
> Ethernet II, Src: de:ad:20:20:22:22 (de:ad:20:20:22:22), Dst: IPv4mcast_7e:fe:fe (01:00:5e:7e:fe:fe)
> Internet Protocol Version 4, Src: 10.10.20.2, Dst: 239.254.254.254
> User Datagram Protocol, Src Port: 4342, Dst Port: 4342
v Locator/ID Separation Protocol
  0100 .... = Type: Map-Notify (4)
  .... 0... = I bit (xTR-ID present): Not set
  .... .0.. = R bit (Built for an RTR): Not set
  .... ..00 0000 0000 0000 = Reserved bits: 0x00000
Record Count: 4
Nonce: 0x0000000000000000
Key ID: 0x0000
Authentication Data Length: 0
Authentication Data: <MISSING>
> Mapping Record 1, EID Prefix: 172.16.144.2/32, TTL: -1610285056, Action: No-Action, Not Authoritative
> Mapping Record 2, EID Prefix: 172.16.144.111/32, TTL: -1610285056, Action: No-Action, Not Authoritative
> Mapping Record 3, EID Prefix: 172.16.144.252/32, TTL: -1610285056, Action: No-Action, Not Authoritative
> Mapping Record 4, EID Prefix: 172.16.144.254/32, TTL: -1610285056, Action: No-Action, Not Authoritative

```

将LISP /32路由重分发到IGP

#这是IGP辅助模式的关键；任何/32 LISP路由都将重分发到IGP;这可通过在EIGRP下应用的“redistribute LISP”命令实现。

#在重分发后，任何/32主机路由都将被视为EIGRP外部路由。调整EIGRP管理距离以使其更高。这是为了确保LISP路由与传入EIGRP外部路由保持在URIB中。例如：DC1-Agg1和DC1-Agg2是具有DC1核心的EIGRP邻居。DC1-AGG1通过重分发将/32路由注入DC1-Core。既然DC1-Core是与DC1-Agg2的EIGRP邻居，那么如果EIGRP AD是170，则同一路由可能返回到DC1-Agg2，并且有机会赢取LISP路由（AD为240）；因此，为避免这种情况，EIGRP外部路由AD已修改为245。

DC1-Agg交换机获知的/32路由将重分发到EIGRP，DC1-core条目如下所示。

```

DC1-CORE# sh ip route 172.16.144.1
IP Route Table for VRF "default"
'*' denotes best ucast next-hop
***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>

172.16.144.1/32, ubest/mbest: 2/0
  *via 192.168.98.1, Eth3/20.111, [170/51456], 00:00:01, eigrp-100, external
  *via 192.168.98.5, Eth3/22.112, [170/51456], 18:14:51, eigrp-100, external

```

#该路由存在于全局路由表中，且核心端未配置VRF。

#由于在AGG交换机上配置了“redistribute direct”，因此核心层还将为父子网提供/24 ECMP路由，如下所示。这将有助于吸引静默主机（没有/32路由）的流量。

```

DC1-CORE# sh ip route 172.16.144.10          # Checking for a non existent Host
172.16.144.10
IP Route Table for VRF "default"
'*' denotes best ucast next-hop
***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>

```

```
172.16.144.0/24, ubest/mbest: 2/0
  *via 192.168.98.1, Eth3/20.111, [170/51456], 00:02:13, eigrp-100, external
  *via 192.168.98.5, Eth3/22.112, [170/51456], 18:17:03, eigrp-100, external
```

#此外，DC1和DC2核心都会看到/24 ECMP路由

```
Branch1-Router# sh ip route 172.16.144.10
Routing entry for 172.16.144.0/24
  Known via "eigrp 100", distance 170, metric 51712, type external
  Redistributing via eigrp 100
  Last update from 192.168.99.2 on GigabitEthernet0/0/1, 00:00:17 ago
  Routing Descriptor Blocks:
    192.168.99.2, from 192.168.99.2, 00:00:17 ago, via GigabitEthernet0/0/1      # 192.168.99.2
  is DC2-Core
    Route metric is 51712, traffic share count is 1
    Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit
    Reliability 255/255, minimum MTU 1492 bytes
    Loading 1/255, Hops 2
  * 192.168.99.1, from 192.168.99.1, 00:00:17 ago, via GigabitEthernet0/0/1      # 192.168.99.1
  is DC1-Core
    Route metric is 51712, traffic share count is 1
    Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit
    Reliability 255/255, minimum MTU 1492 bytes
    Loading 1/255, Hops 2
```

#此路由可确保分支主机可以连接到位于任一位置的静默主机。

DC内VLAN间数据包的路径

#当DC1-Host1 -> 172.16.144.1尝试到达DC2-Host1-> 172.16.144.2时，这是数据中心间VLAN流量。DC1-Host 1发出ARP请求，该请求将一路穿过OTV并到达DC2-Host1

DC2-Host1以返回DC1-Host1的ARP应答作出响应

#后续ICMP数据包通过OTV发送

DC间VLAN间的数据包路径 (从VLAN 144到VLAN 244)

#当DC1-Host1-> 172.16.144.1尝试到达DC2-Host2-> 172.16.244.2时，数据包不会从DC1的vlan 144路由到244;相反，它遵循从DC1-Agg到DC1-Core的路由路径，然后到达DC2-Core，最终路由将由DC2-Agg交换机完成到目的Vlan-244。

#从DC1-Host1到DC2-Host2的traceroute如下所示。

```
DC1-HOST# traceroute 172.16.244.2 vrf vlan144
traceroute to 172.16.244.2 (172.16.244.2), 30 hops max, 40 byte packets
 1 172.16.144.250 (172.16.144.250) 1.149 ms 0.841 ms 0.866 ms
# DC1-AGG1
 2 192.168.98.2 (192.168.98.2) 1.004 ms 0.67 ms 0.669 ms
# DC1-CORE
 3 192.168.99.2 (192.168.99.2) 0.756 ms 0.727 ms 0.714 ms
# DC2-CORE
 4 192.168.94.5 (192.168.94.5) 1.041 ms 0.937 ms 192.168.94.1 (192.168.94.1) 1.144 ms
# DC2-Agg1/DC2-Agg2
 5 172.16.244.2 (172.16.244.2) 2.314 ms * 2.046 ms
# DC2-Host2
```

DC间VLAN间数据包的路径 (从VRF租户-1到VRF租户-2)

#这将遵循与从一个VLAN到另一个VLAN的DC间通信相同的方法 (上例)

#当DC1-host1-> 172.16.144.1尝试到达DC2-Host3-> 172.16.145.2时，这是发往Vlan 144 (VRF租户-1) 的DC间流量，发往Vlan 145 (VRF租户-2)。与常规N7k OTV部署不同，此流量的处理方式略有不同。DC1端不会发生任何VLAN间路由；相反，此流量将被路由并发送到DC1核心，而核心层将通过IGP进一步将其路由到DC2核心层

#为了本文档，核心交换机按站点执行VRF间泄漏。它可以是任何设备 (如防火墙) ；如果VRF间泄漏存在或不存在，则从LISP配置角度看没有更改。

```
DC1-AGG1# sh ip route 172.16.145.2 vrf tenant-1
IP Route Table for VRF "tenant-1"
 '*' denotes best ucast next-hop
 '**' denotes best mcast next-hop
 '[x/y]' denotes [preference/metric]
 '%<string>' in via output denotes VRF <string>
```

```
172.16.145.2/32, ubest/mbest: 1/0
```

```
 *via 192.168.98.2, Eth3/6.111, [245/51968], 00:00:46, eigrp-100, external
```

#从DC1-Host1到DC2-Host3的Traceroute将显示与其未通过VLAN间路由的路由相同，而是通过核心路由的第3层。简而言之，VLAN间流量将不使用OTV。

```
DC1-HOST# traceroute 172.16.145.2 vrf vlan144
traceroute to 172.16.145.2 (172.16.145.2), 30 hops max, 40 byte packets
 1 172.16.144.250 (172.16.144.250) 1.049 ms 0.811 ms 0.81 ms #
DC1-AGG1
 2 192.168.98.2 (192.168.98.2) 0.844 ms 0.692 ms 0.686 ms #
DC1-CORE
 3 192.168.99.2 (192.168.99.2) 0.814 ms 0.712 ms 0.735 ms #
DC2-CORE
 4 192.168.194.1 (192.168.194.1) 0.893 ms 0.759 ms 192.168.194.5 (192.168.194.5) 0.89 ms #
DC2-Agg1/DC2-Agg2
 5 172.16.145.2 (172.16.145.2) 1.288 ms * 1.98 ms #
DC2-Host3
DC1-HOST#
```

当Branch-1主机尝试到达DC2中存在的静默主机时数据包的路径

Branch-1-172.17.200.1中的主机尝试到达DC2-Silent Host- 172.16.144.119。由于主机处于静默状态，因此DC2中不会存在任何/32路由。

```
DC2-AGG1# show ip route 172.16.144.119 vr tenant-1
IP Route Table for VRF "tenant-1"
 '*' denotes best ucast next-hop
 '**' denotes best mcast next-hop
 '[x/y]' denotes [preference/metric]
 '%<string>' in via output denotes VRF <string>
```

```
172.16.144.0/25, ubest/mbest: 1/0
```

```
 *via Null0, [240/1], 20:48:29, lisp, dyn-eid
```

```
DC2-AGG2# show ip route 172.16.144.119 vr tenant-1
IP Route Table for VRF "tenant-1"
```

'*' denotes best ucast next-hop
'**' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>

```
172.16.144.0/25, ubest/mbest: 1/0  
  *via Null0, [240/1], 20:48:13, lisp, dyn-eid
```

#根据LISP设计，路由172.16.144.119将与172.16.144.0/25 null0路由匹配。

#当Branch路由器收到目的IP = 172.16.144.119的数据包时，URIB具有到DC1核和DC2核的ECMP /24路由。这实质上意味着数据包将发送到其中一台核心层交换机。

```
Branch1-Router# sh ip route 172.16.144.119  
Routing entry for 172.16.144.0/24  
  Known via "eigrp 100", distance 170, metric 51712, type external  
  Redistributing via eigrp 100  
  Last update from 192.168.99.2 on GigabitEthernet0/0/1, 00:08:54 ago  
  Routing Descriptor Blocks:  
    192.168.99.2, from 192.168.99.2, 00:08:54 ago, via GigabitEthernet0/0/1  
      Route metric is 51712, traffic share count is 1  
      Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit  
      Reliability 255/255, minimum MTU 1492 bytes  
      Loading 1/255, Hops 2  
  * 192.168.99.1, from 192.168.99.1, 00:08:54 ago, via GigabitEthernet0/0/1  
      Route metric is 51712, traffic share count is 1  
      Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit  
      Reliability 255/255, minimum MTU 1492 bytes  
      Loading 1/255, Hops 2
```

```
Branch1-Router#sh ip cef exact-route 172.17.200.1 172.16.144.119 dest-port 1  
172.17.200.1 -> 172.16.144.119 =>IP adj out of GigabitEthernet0/0/1, addr 192.168.99.1
```

#根据CEF，数据包散列到192.168.99.1（即DC1-Core）

DC1-Core有2条ECMP路径；一个指向DC1-Agg1（HSRP活动），另一个指向DC1-Agg2（HSRP备用）。从路由散列中，选择的路径将是DC1-Agg2。

```
DC1-CORE# sh routing hash 172.17.200.1 172.16.144.119 1 1 Load-share parameters used for  
software forwarding: load-share mode: address source-destination port source-destination  
Universal-id seed: 0xfdba3ebe Hash for VRF "default" Hash Type is 1 Hashing to path  
*192.168.98.5 Eth3/22.112  
For route:  
172.16.144.0/24, ubest/mbest: 2/0  
  *via 192.168.98.1, Eth3/20.111, [170/51456], 00:19:57, eigrp-100, external  
  *via 192.168.98.5, Eth3/22.112, [170/51456], 18:34:47, eigrp-100, external
```

```
DC1-CORE# sh cdp nei int e3/22  
Capability Codes: R - Router, T - Trans-Bridge, B - Source-Route-Bridge  
  S - Switch, H - Host, I - IGMP, r - Repeater,  
  V - VoIP-Phone, D - Remotely-Managed-Device,  
  s - Supports-STP-Dispute
```

Device-ID	Local Interface	Hldtme	Capability	Platform	Port ID
DC1-AGG2(JAF1534CHCJ)	Eth3/22	172	R S s	N7K-C7009	Eth3/7

#由于DC1-Agg2在URIB中没有任何条目，因此数据包将被收集并发送到CPU，这将强制DC1-Agg2从SVI IP地址生成ARP请求，如下所示。

```
2020-02-18 15:09:05.673165 172.17.200.1 -> 172.16.144.119 ICMP 114 Echo (ping) request
id=0x0022, seq=0/0, ttl=254
2020-02-18 15:09:05.675041 de:ad:20:19:22:22 -> Broadcast ARP 60 Who has 172.16.144.119? Tell
172.16.144.251
```

#此ARP请求是广播，它在整个第2层域内传播，该域还包括通过OTV扩展的DC2。

DC2-Silent Host现在响应来自DC1-Agg2的ARP请求

DC1-Agg2从静默主机收到此ARP应答

```
2020-02-18 15:09:05.675797 64:12:25:97:46:41 -> de:ad:20:19:22:22 ARP 60 172.16.144.119 is at
64:12:25:97:46:41
```

#由于收到的数据包是ARP（用作LISP的提示），因此会从HSRP VIP-> 172.16.144.254生成ICMP ECHO请求，该请求发往静默主机> 172.16.144.119。从HSRP VIP获取数据包的目的是了解主机是本地还是远程。如果主机是远程的，则FHRP活动也存在于远程数据中心中，该数据中心将从主机捕获ICMP ECHO应答数据包，因此这会导致DC2-Agg2（即HSRP活动）了解此条目，LISP进程现在将基于此IP数据包创建EID学习。最初从HSRP VIP发出ICMP ECHO请求的DC1-Agg2从未收到响应，因此DC1端永远不会发生终端学习；而是在DC2侧。

```
DC2-AGG2# show lisp dynamic-eid detail vrf tenant-1
LISP Dynamic EID Information for VRF "tenant-1"
Dynamic-EID name: VLAN144
  Database-mapping [2] EID-prefix: 172.16.144.0/24, LSBs: 0x00000003
    Locator: 10.10.20.1, priority: 50, weight: 50
      Uptime: 21:50:32, state: up
    Locator: 10.10.20.2, priority: 50, weight: 50
      Uptime: 21:50:13, state: up, local
Registering more-specific dynamic-EIDs
Registering routes: disabled
Allowed-list filter: none applied
Map-Server(s): none configured, use global Map-Server
Site-based multicast Map-Notify group: 239.254.254.254
Extended Subnet Mode configured on 1 interfaces
Number of roaming dynamic-EIDs discovered: 3
Last dynamic-EID discovered: 172.16.144.254, 00:01:10 ago
Roaming dynamic-EIDs:
  172.16.144.2, Vlan144, uptime: 19:09:07, last activity: 00:05:21
    Discovered by: packet reception
  172.16.144.119, Vlan144, uptime: 00:05:55, last activity: 00:05:55 Discovered by: packet
reception
  172.16.144.252, Vlan144, uptime: 3d21h, last activity: 00:01:10
    Discovered by: packet reception
```

Secure-handoff pending for sources: none

#一旦LISP进程知道DC2-Agg2（HSRP活动）上的EID，它将

a)在本地安装/32

b)重分布到DC2-Core的路由

c)在Vlan中以组播消息的形式发送基于站点的通知（在本例中，消息将发往组 —> 239.254.254.254）

```

DC2-AGG1# show lisp dynamic-eid detail vrf tenant-1
LISP Dynamic EID Information for VRF "tenant-1"
Dynamic-EID name: VLAN144
  Database-mapping [2] EID-prefix: 172.16.144.0/24, LSBs: 0x00000003
    Locator: 10.10.20.1, priority: 50, weight: 50
      Uptime: 21:52:39, state: up, local
    Locator: 10.10.20.2, priority: 50, weight: 50
      Uptime: 21:52:08, state: up
  Registering more-specific dynamic-EIDs
  Registering routes: disabled
  Allowed-list filter: none applied
  Map-Server(s): none configured, use global Map-Server
  Site-based multicast Map-Notify group: 239.254.254.254
  Extended Subnet Mode configured on 1 interfaces
  Number of roaming dynamic-EIDs discovered: 4
  Last dynamic-EID discovered: 172.16.144.254, 00:03:07 ago
  Roaming dynamic-EIDs:
    172.16.144.2, Vlan144, uptime: 19:11:04, last activity: 00:00:21
      Discovered by: site-based Map-Notify
    172.16.144.110, Vlan144, uptime: 20:04:09, last activity: 20:04:09
      Discovered by: site-based Map-Notify
    172.16.144.119, Vlan144, uptime: 00:07:52, last activity: 00:00:21 Discovered by: site-based
Map-Notify
    172.16.144.252, Vlan144, uptime: 21:50:51, last activity: 00:00:21
      Discovered by: site-based Map-Notify

```

Secure-handoff pending for sources: none

#最后，Branch-router1将收到此/32路由，这将导致Branch路由器将流量发送到正确的DC2核心交换机。

```

Branch1-Router# sh ip route 172.16.144.119
Routing entry for 172.16.144.119/32
  Known via "eigrp 100", distance 170, metric 51712, type external
  Redistributing via eigrp 100
  Last update from 192.168.99.2 on GigabitEthernet0/0/1, 00:06:25 ago
  Routing Descriptor Blocks:
  * 192.168.99.2, from 192.168.99.2, 00:06:25 ago, via GigabitEthernet0/0/1
    Route metric is 51712, traffic share count is 1
    Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit
    Reliability 255/255, minimum MTU 1492 bytes
    Loading 1/255, Hops 2

```

主机从DC1移动（漫游）到DC2时的事件顺序

#考虑到此拓扑上配置了L2扩展，主机可以从DC1移动到DC2。

Host-> 172.16.144.100最初在VLAN 144和DC1中。

#当主机在DC1上线时，DC1-Agg1和DC1-Agg2交换机内的路由将如下所示

```

DC1-AGG1# sh ip route 172.16.144.100 vrf tenant-1
IP Route Table for VRF "tenant-1"
 '*' denotes best ucast next-hop
 *** denotes best mcast next-hop
 '[x/y]' denotes [preference/metric]
 '%<string>' in via output denotes VRF <string>

```

```
172.16.144.100/32, ubest/mbest: 1/0, attached
  *via 172.16.144.100, Vlan144, [240/1], 00:05:03, lisp, dyn-eid
  via 172.16.144.100, Vlan144, [250/0], 00:05:05, am
```

```
DC1-AGG2# sh ip route 172.16.144.100 vrf tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
 '**' denotes best mcast next-hop
 '[x/y]' denotes [preference/metric]
 '%<string>' in via output denotes VRF <string>
```

```
172.16.144.100/32, ubest/mbest: 1/0, attached
  *via 172.16.144.100, Vlan144, [240/1], 00:08:05, lisp, dyn-eid
  via 172.16.144.100, Vlan144, [250/0], 00:08:07, am
```

#分支路由器将具有指向DC1-Core的路由，如下所示，而traceroute将指向DC1 Core/agg交换机以到达DC1中的主机

```
Branch1-Router#sh ip route 172.16.144.100
Routing entry for 172.16.144.100/32
  Known via "eigrp 100", distance 170, metric 51712, type external
  Redistributing via eigrp 100
  Last update from 192.168.99.1 on GigabitEthernet0/0/1, 00:00:06 ago
  Routing Descriptor Blocks:
  * 192.168.99.1, from 192.168.99.1, 00:00:06 ago, via GigabitEthernet0/0/1
    Route metric is 51712, traffic share count is 1
    Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit
    Reliability 255/255, minimum MTU 1492 bytes
    Loading 1/255, Hops 2
```

```
Branch1-Router#traceroute 172.16.144.100 source 172.17.200.1
Type escape sequence to abort.
Tracing the route to 172.16.144.100
VRF info: (vrf in name/id, vrf out name/id)
  1 192.168.99.1 1 msec 1 msec 0 msec          # DC1-Core
  2 192.168.98.5 1 msec 1 msec                # DC1-Agg2
    192.168.98.1 1 msec                       # DC1-Agg1
  3 172.16.144.100 1 msec 0 msec 1 msec       # DC1-Host
```

#当主机移动到DC2时，它会在Vlan 144中发送GARP。在DC2-Agg交换机上会看到此情况

```
2020-02-24 22:23:05.024902 Cisco_5a:4a:e7 -> Broadcast ARP 60 Gratuitous ARP for
172.16.144.100 (Request)
```

#一旦收到包含ARP/GARP/RARP的数据包，就会触发本地化机制，向源自VIP的主机发出ICMP回应请求

```
2020-02-24 22:23:05.026781 172.16.144.254 -> 172.16.144.100 ICMP 60 Echo (ping) request
id=0xac10, seq=0/0, ttl=128
```

#主机172.16.144.100现在将响应HSRP VIP

```
2020-02-24 22:23:07.035292 172.16.144.100 -> 172.16.144.254 ICMP 60 Echo (ping) reply
id=0xac10, seq=0/0, ttl=255
```

#一旦在DC2-Agg1处收到IP数据包，这将导致LISP检测EID并在主机的路由表中输入条目，并启动到EIGRP的重分发过程

```
DC2-AGG1# sh ip route 172.16.144.100 vrf tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
'***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>
```

```
172.16.144.100/32, ubest/mbest: 1/0, attached
  *via 172.16.144.100, Vlan144, [240/1], 00:00:30, lisp, dyn-eid
  via 172.16.144.100, Vlan144, [250/0], 00:00:32, am
```

#当重分发到位后，DC1-agg站点（该站点是此主机的原始所有者）现在将看到指向EIGRP的RIB中的更改

```
DC1-AGG1# sh ip route 172.16.144.100 vrf tenant-1
IP Route Table for VRF "tenant-1"
'*' denotes best ucast next-hop
'***' denotes best mcast next-hop
'[x/y]' denotes [preference/metric]
'%<string>' in via output denotes VRF <string>
```

```
172.16.144.100/32, ubest/mbest: 1/0
  *via 192.168.98.2, Eth3/6.111, [245/51968], 00:03:47, eigrp-100, external
```

#远程分支路由器现在将看到路由更改，traceroutes将反映到DC2核心/Agg交换机的路径更改，如下所示

```
Branch1-Router#sh ip route 172.16.144.100
Routing entry for 172.16.144.100/32
  Known via "eigrp 100", distance 170, metric 51712, type external
  Redistributing via eigrp 100
  Last update from 192.168.99.2 on GigabitEthernet0/0/1, 00:00:00 ago
  Routing Descriptor Blocks:
  * 192.168.99.2, from 192.168.99.2, 00:00:00 ago, via GigabitEthernet0/0/1
    Route metric is 51712, traffic share count is 1
    Total delay is 1020 microseconds, minimum bandwidth is 100000 Kbit
    Reliability 255/255, minimum MTU 1492 bytes
    Loading 1/255, Hops 2
```

```
Branch1-Router#traceroute 172.16.144.100 source 172.17.200.1
Type escape sequence to abort.
Tracing the route to 172.16.144.100
VRF info: (vrf in name/id, vrf out name/id)
  1 192.168.99.2 1 msec 0 msec 1 msec          # DC2-Core
  2 192.168.94.1 1 msec 1 msec 1 msec        # DC2-Agg1
  3 172.16.144.100 0 msec 0 msec 1 msec      # Host-after move to DC2
```

有用的验证命令

show lisp dynamic-eid detail vrf <VRF Name>

Show ip route lisp vrf <VRF Name>

show lisp dynamic-eid summary vrf <VRF Name>