BGP RR Scale Considerations and KPI Monitoring

Contents

Introduction

This document describes the major contributors to the maximum scale a Border Gateway Protocol (BGP) Route-Reflectors (RR) can achieve and guidance on BGP RR performance monitoring.

HW/SW Platform Selection

A high-scale BGP RR is typically not in the forwarding path of packets carrying services provided by an Internet Service Provider. Therefore, the hardware requirements for a BGP RR and routers that are predominantly forwarding packets in the data path are different. Standard routers are built with a powerful data-path forwarding element and a comparatively moderate control-path element. A BGP RR performs all its tasks in a control plan.

Within the Cisco IOS® XR family of products, you can choose between 3 flavours of HW/SW platforms for a BGP RR role:

As of this writing, XRv9k Appliance is the optimal platform choice for BGP RR because it provides for highest control plane capacity with maximum performance.

Scale And Performance Considerations

The supported scale of data-plane entities is relatively easy to express because the performance of the datapath element seldom depends on the scale. For example, a TCAM lookup takes the same time regardless of the number of active TCAM entries.

The supported scale of control-plane entities is often much more complex because the scale and performance are interconnected. Consider a BGP RR with 1M routes. The work that a BGP process must carry out to maintain this BGP table depends on:

- 1. How many BGP peers are active?
- 2. What address families are active?
- 3. How are they distributed into update groups?
- 4. The complexity of RPLs (Route Policies)
- 5. Frequency of updates (incoming updates and also outgoing updates advertisement interval).
- 6. TCP MSS, Interface/Path MTU tuning this will aid in better performance
- 7. If dual-RP, is NSR enabled
- 8. Any known slow-peers, that aren't in separate update-group
- 9. Nexthop trigger-delay value

Number Of BGP Peers

The number of BGP peers is usually the first and unfortunately, often the only thing that comes to mind when considering the BGP scale. While the supported BGP scale cannot be represented without mentioning the number of BGP peers, it is not the most important factor. Many other aspects are equally relevant.

Address Families

The type of address family (AF) is an important factor in BGP performance considerations because in typical deployments it impacts the size of a single route. The number of IPv4 routes that can be packed into a single TCP segment is significantly higher than the number of VPNv4 routes. Hence for the same scale of BGP table changes, an IPv4 BGP RR has less work to do compared to a VPNv4 BGP RR. Obviously, in deployments where a significant number of communities is added to every route, the difference between AFs becomes less significant, but the size of a single route is then even bigger and requires consideration.

Number Of Update Groups

The BGP process prepares a single update for all members of the same update group. Then TCP process splits the update data into a required number of TCP segments (depending on TCP MSS) towards each member of the update group. You can see the active update groups and their members by using the show bgp update-group command. You can influence which and how many peers are members of an update group by creating a common outbound policy for a group of peers you want to be in the same update group. A single update sent by the BGP RR to a high number of BGP RR clients can trigger a burst of TCP ACKs that can be dropped in Local Packet Transport Service (LPTS) component of Cisco IOS XR routers.

Complexity of RPLs (Route Policies)

The complexity of route policies used by BGP impacts the BGP process performance. Every received or sent route must be evaluated against the configured route policy. A very long policy requires many CPU cycles to be spent on this action. A route policy that includes a regular expression is especially heavy on processing. A regular expression helps you express the route policy in a lesser number of lines but requires more CPU cycles while processing than the equivalent route policy that does not use regular expression.

Frequency Of Updates

The frequency of updates has an important bearing on the BGP scale. The number of updates is often difficult to predict. You can influence the frequency of updates by using the "**advertisementinterval**" command, which sets the minimum interval between the sending of BGP) routing updates. The default value towards iBGP peers is 0 seconds and 30 towards eBGP peers it is 30 seconds.

TCP MSS And Interface/Path MTU

Splitting an update into many TCP segments can put a lot of strain on TCP process resources in a high-scale and high-update frequency environment. A bigger path MTU and bigger TCP MSS are better for BGP and TCP performance.

NSR On Dual-RP Routers

NSR is a great feature for redundancy, but it does have an impact on BGP performance. On Cisco IOS XR routers both RPs are receiving simultaneously every BGP update directly from the NPU on the ingress line card, which means that the Active RP does not have to spend time on replicating the update to the Standby RP. However, every update generated by the Active RP must be sent to the Standby RP and from there to the BGP peer. This allows the Standby RP to be always up to date on the sequence and acknowledgement numbers but does have an impact on the overall BGP performance. This is why it is recommended that a BGP RR is a single-RP router.

Slow Peers

A slow peer can slow down the updates towards all members of the update group because the BGP process must keep the update in its memory until all peers have acknowledged it. If you know that some peers are much slower (for example routers in a legacy part of the network), separate them upfront into an update group. By default, Cisco IOS XR reports a slow peer via syslog message. You can create static slow peers (that never share the update group with others) or fine-tune the dynamic slow peer behaviour by using the BGP slow-peer configuration command in global or per-neighbour config mode. A good further read on this can be found in [Troubleshoot Slow BGP Convergence Due to Suboptimal Route Policies on IOS-XR](https://xrdocs.io/asr9k/tutorials/troubleshoot-slow-bgp-convergence-due-to-rpl/) on the Cisco xrdocs.io portal.

Nexthop trigger-delay

If multiple BGP next-hops change in a short time interval and the critical nexthop trigger-delay value of zero is configured in an address family (AF) with a high number of routes, a full walk of the AF must be executed on every next-hop change event. Repeated walks of that AF increase the convergence time in address families with lower critical nexthop trigger-delay values. You can see the next-hop trigger-delay values by running the "show bgp all all nexthops" command.

Example Of Validated Multidimensional BGP RR Scale

Multidimensional scale results, especially for control plane features, are highly dependent on the specific test environment. Test results can vary significantly if some of the parameters are changed.

Design Considerations

There are two approaches to BGP RR placement in the network:

- 1. Centralised/Flat BGP RR design.
- 2. Distributed/Hierarchical BGP RR design.

In a centralised/flat design, all BGP RR clients in the network establish BGP peering with a set (usually a pair) of BGP RR devices that hold the exact same information. This approach is simple to implement and works well in small to moderate scale networks. Any change in BGP table is propagated quickly to all BGP RR clients. As the number of BGP RR clients grows, the design can reach a scale limit when the number of TCP connections on the BGP RR devices grows to the extent where their performance is impacted.

In a distributed/hierarchical design, network is split into several regions. All routers in a region establish BGP peering with a set (usually a pair) of BGP RR devices that hold the exact same information. These BGP RR devices act as BGP RR clients to another set (usually a pair) of BGP RR devices. This design approach allows for easy network expansion, while keeping the number of TCP connections on every single BGP RR under certain limit.

Another design consideration is tailoring the scope of recipients of BGP updates. Depending on the VRF distribution among BGP RR clients, it's worth considering the RT Constrained Route Distribution. If all BGP RR clients have interfaces in the same VRF, RT Constrained Route Distribution does not bring many benefits. However, if VRFs are sparsely distributed among all BGP RR clients, use of RT Constrained Route Distribution significantly reduces f the load on the bgp process on the BGP RR.

Monitor BGP Key Performance Indicators (KPI)

Monitoring of BGP RR's Key Performance Indicators (KPI) is important for ensuring proper network operation.

A significant change in the network topology (for example a major DWDM link flap) can trigger routing updates that generate excessive traffic towards and/or from the BGP RR. Significant traffic hitting the BGP RR typically carries:

- 1. Updates from BGP peers.
- 2. TCP ACKs generated by the BGP peers, in response to updates sent by BGP RR & vice-versa

This section of the document explains the KPI that need to be monitored on a typical BGP RR and also how to tell which of the two significant BGP traffic types are causing high control plane traffic rate.

Path of BGP packets inside the router can be depicted as follows:

Punt

Ethernet controller -(packet)-> datapath forwarder -(packet)-> LPTS -(packet)-> SPP -(packet) -> NetIO -(packet)-> TCP -(message)-> BGP

Inject

BGP -(message)-> TCP -(packet)-> NetIO -(packet)-> SPP -(packet) -> datapath forwarder -(packet)- > Ethernet controller

KPIs can be split into:

Essentials:

- 1. Datapath Forwarder
- 2. LPTS (hardware punt policers settings, accept counters and drop counters)
- 3. SPP
- 4. NetIO
- 5. IPC queues (NetIO $\leq \Rightarrow$ TCP $\leq \Rightarrow$ BGP)
- 6. BGP InQ/OutQ sizes

Optional:

- 1. CPU utilisation
- 2. Memory utilisation
- 3. TCP statistics
- 4. BGP process performance
- 5. BGP convergence

Monitor Datapath Forwarder

On XRv9000 the datapath forwarder is the Data Plane Agent (DPA), while on ASR9000 platforms it is the Network Processor (NP).

Monitor XRv9000 Data Plane Agent (DPA)

Useful command to see the load and statistics of the DPA is:

```
   show controllers dpa statistics global
```
This command shows all the non-zero counter, that give you insight into the type and number of packets punted from network interfaces to RP CPU, injected from RP CPU towards network interfaces, and the numer of packets dropped:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show controllers dpa statistics global

Monitor ASR9000 Network Processor (NP)

Useful commands to see the load and statistics of each NP in the system are:

show controllers np load all

show controllers np counters all

NP on ASR9000 has a rich set of counters that show you the number, rate and type of processed and dropped packets,.

<#root>

RP/0/RSP0/CPU0:ASR9k-B#

show controllers np load all

Node: 0/0/CPU0:

-- Load Packet Rate NP0: 0% utilization 937 pps NP1: 0% utilization 538 pps RP/0/RSP0/CPU0:ASR9k-B#

RP/0/RSP0/CPU0:ASR9k-B#

show controllers np counters all

Node: 0/0/CPU0:


```
Monitor LPTS
```
As a standard BGP RR is not in the forwarding path, all packets received on network interface are punted to control-plane. The data-path element on a BGP RR performs a small number of simple operations before packets are punted to control-plane. As the data-path element is unlikely to be a point of congestion, the only element on the line card that needs monitoring are the LPTS stats.

Please note that in case of XRv9k, hardware statisitics map to the vPP

Command:

show lpts pifib hardware police location <location> | inc "Node|flow_type|BGP"

Example:

What to look for: If a significant jump in AggDrops against BGP-known flow type is observed, start looking for network topology changes that have triggered such massive control plane churn.

Telemetry data path:

Cisco-IOS-XR-lpts-pre-ifib-oper:lpts-pifib

Note: LPTS stats counters can be cleared. Your monitoring system must account for this possibility.

Monitor SPP

SPP is the first enttity on the route processor or line card CPU that receives the packet punted from the NP or DPA via internal fabric, and the last point in the software packet processing before it is handed over to the fabric for injection into the NP or DPA.

Relevant commands for SPP monitoring:

show spp node-counters

show spp client

The **show spp node-counters** command shows the rate of punted/injected packets and is easy to read and understand. For BGP sessions the relevant coounters are under **client/pun**t and **client/inject** on the active RP.

The **show spp client** is more rich in output and gives a more detailed insight into number of packets enqueued/dropped towards clients, as well as the high watermark.

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show spp node-counters

0/RP0/CPU0:

socket/rx

0/0/CPU0:

 $\langle \cdot \rangle$. $\langle \cdot \rangle$

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show spp client

Sat Apr 20 17:11:40.725 UTC 0/RP0/CPU0: Clients =======

 $\langle \cdot, \cdot, \cdot \rangle$ netio, JID 254 (pid 4591) -- Reconnect Pending: F, Exited: F, Keep Queues: F, Pakman Client: T Quota Current: 0, Limit: 16384, Drops 0 Queues: Key Cur Max Enqueues High Watermark (Timestamp) Drops 0xffffffff 0 10 0 0 (22:13:52.195 Feb 21 24 UTC) 0 0x03000006 0 2048 0 0 (22:13:52.196 Feb 21 24 UTC) 0 0x03000007 0 3072 414881 1 (23:03:30.721 Feb 21 24 UTC) 0 0x03000008 0 1024 5 1 (13:41:28.389 Mar 13 24 UTC) 0 0x03000009 0 2048 180411 1 (23:03:31.565 Feb 21 24 UTC) 0

Monitor NetIO

While the LPTS policer only shows the count of packets accepted or dropped by a corresponding policer, at NetIO level we can see the rate of packets punted to RP CPU. Since on a typical BGP RR the great majority of received packets are BGP packets, the overall NetIO rate indicates very closely the rate of received BGP packets.

<#root>

Command:

show netio rates

Example:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show netio rates

Netio packet rate for node 0/RP0/CPU0

RP/0/RP0/CPU0:xrv9k-01#

What to look for:

• If a significant jump in NetIO rate is observed, start looking for network topology changes that have triggered such massive control plane churn.

Telemetry data path:

• not applicable as telemetry must stream counter values, not rates. BGP-known LPTS policer accept counter can be used on the Telemetry collector to approximate the average rate of received BGP packets from known peers.

Monitor XIPC Queues

On the punt path packets received by NetIO from LPTS are passed on to TCP and BGP. It's important to monitor these queues:

- 1. TCP high priority queue though which NetIO delivers packets to TCP
- 2. BGP control queue
- 3. BGP data queue

On the inject path packets are created by TCP and passed onto NetIO. It's important to monitor these queues:

1. OutputL XIPC queue

Commands:

show netio clients show processes bgp | i "Job Id" show xipcq jid

bgp_job_id>

show xipcq jid

show a queue-id <n>

Examples:

NetIO to TCP, view from NetIO standpoint:

RP/0/RP0/CPU0:xrv9k-01#show netio clients

TCP to NetIO, view from NetIO standpoint::

NetIO to TCP, view from TCP process standpoint:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show processes tcp

| i "Job Id"

 Job Id: 430 RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

430

Mon Apr 17 16:16:11.315 CEST

TCP to BGP:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show processes bgp

| i "Job Id"

 Job Id: 1078 RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

 1078 Mon Apr 17 16:09:33.046 CEST


```
BGP data queue:
```
<#root>

RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

1078

queue-id 1

XIPC_xipcq_12_0_9854_6506_inst_1_data_toapp

```
:
```


BGP control queue:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

1078

queue-id

2

What to look for:

- there must be no drops in relevant queues
- in XIPC queue stats High watermark (HWM) must not exceed the 50% of queue size

For better tracking of high watermark value evolution, you must clear the high watermark value after every read. Note that this does not clear only the HWM counter but also clears all queue statistics. The format of the command for clearing the XIPC queue statistics is: clear xipcq statistics queue-name <queue_name>

As the queue name often includes the Process ID (PID), the queue name changes after process restart. Some examples of commands for clearing the relevant queues statistics:

```
  clear xipcq statistics queue-name XIPC_tcp_i0
  clear xipcq statistics queue-name XIPC_tcp_i1
  clear xipcq statistics queue-name XIPC_xipcq_12_0_9854_6506_inst_1_data_toapp
  clear xipcq statistics queue-name XIPC_xipcq_12_0_9854_6506_inst_1_ctrl_toapp
```
Telemetry path:

• There are no Telemetry sensor paths for XIPC.

Monitor BGP Input And Output Queues

BGP maintains an input and output queue for every BGP peer. Data sits in InQ when TCP has passed it to BGP, but BGP hasn't processed them yet. Data sits in OutQ while BGP waits on TCP to split the data into packets and transmit them. The instantaneous size of BGP InQ/OutQ provides a good indication of how busy the BGP process is.

Command:

show bgp <AFI> <SAFI> summary

Example:

RP/0/RP0/CPU0:xrv9k-01#show bgp all all summary Address Family: VPNv4 Unicast -----------------------------

BGP router identifier 192.168.0.1, local AS number 65000 BGP generic scan interval 60 secs BGP table state: Active Table ID: 0x0 BGP main routing table version 2208096 BGP scan interval 60 secs

BGP is operating in STANDALONE mode.

What to look for:

- The size of InQ/OutQ must be zero when the network is stable. It changes quickly when updates are exchanged.
- InQ/OutQ size must not monotonously increase over time.

Telemetry path:

• Cisco-IOS-XR-ipv4-bgp-oper:bgp

Monitor BGP Message Rates

Some BGP neighbors can continuously send updates or withdrawals if network topology is unstable. The BGP RR must then replicate such routing table change thousands of times to all its RR clients. Therefore it is important to monitor the message rates received from neighbors, to track sources of instabilities.

Command:

show bgp <AFI> <SAFI> summary

Example:

RP/0/RP0/CPU0:xrv9k-01#show bgp all all summary Address Family: VPNv4 Unicast -----------------------------

BGP router identifier 192.168.0.1, local AS number 65000 BGP generic scan interval 60 secs BGP table state: Active Table ID: 0x0 BGP main routing table version 2208096 BGP scan interval 60 secs

BGP is operating in STANDALONE mode.

RR clients queues have roughly the same amount of MsgSent but some neighbors can have a number of MsgRcvd higher than others. You must capture multiple snapshots of this command to assess the rate of messages.

Once you have identified the offending peers, you can then go through other commands like **show bgp neighbor <neighbor> detail** and **show bgp neighbor <neighbor> performance-statistics** or **show bgp recent-prefixe**s to try to understand which prefixes are flapping and whether it is always the same ones or different ones.

running a command like show bgp all all summary , you see the same counters per neighbor in the sections for the various address-families. They do not represent the number of messages received/sent from/to that neighbor for that address-family but across address-families.

Monitor CPU Utilisation

CPU utilisation must be monitored on every router, but on a router with a high number of CPU cores dedicated to control plane some readouts can be unintuitive. On an BGP RR with a high number of CPU cores dedicated to Routing Processor (RP), as in the case of XRv9k Appliance, active threads run on different CPU cores, while a number of CPU cores remains idle. As a consequence some CPU cores can be very busy, but the overall CPU utilisation calculated across all CPU cores remains moderate.

Therefore, for proper monitoring of CPU cores utilisation via CLI, use the **show processes cpu thread** command.

Monitor TCP Statistics

Cisco IOS® maintains detailed statistics on every TCP session. CLI command **show tcp brief** displays the list of all existing TCP sessions. In this summary output, for every TCP session you can see this information:

- **PCB:** unique TCP session identifier.
- **VRF-ID**: the ID of the VRF in which the session exists.
	- To see the correspondng VRF name, run this command:
	- show cef vrf all summary | utility egrep "^VRF:|Vrfid" | utility egrep -B1 <VRF-ID>
- **Recv-Q**: instantaneous size of the receive Q. Receive queue holds packets received from NetIO. The **tcp** process extracts the data from a packet and sends it to the corresponding application.
- Send-Q: instantaneous size of the send Q. Send queue holds data received from an application. The **tcp** process splits the data into TCP segments (dictated by negotiated maximum segment size - TCP MSS), encapsulates every segment into a layer 3 header of corresponding address family (IPv4 or IPv6) and sends the packet to NetIO.
- Local Address: local IPv4 or IPv6 address associated with the TCP socket. TCP sessions in LISTEN state are typically bound to "**any**" IP address, which is represented as "0.0.0.0" or "::" in case of IPv4 or IPv6 respectively.
- Foreign Address: remote IPv4 or IPv6 address associated with the TCP socket. TCP sessions in LISTEN state are typically bound to "**any**" IP address, which is represented as "0.0.0.0" or "::" in case of IPv4 or IPv6 respectively.
- **State**: TCP session state. Possible TCP session states are: LISTEN, SYNSENT, SYNRCVD, ESTAB, LASTACK, CLOSING, CLOSEWAIT, FINWAIT1, FINWAIT2, TIMEWAIT, CLOSED.

As the well-known BGP port number is 179, you can limit the displayed TCP sessions to those that are associated with the BGP application.

Example:

You can use the displayed PCB value to obtain the statistics for a particular TCP session. CLI commands that provide insight into TCP process statistics:

Global:

show tcp statistics clients location <active RP>

show tcp statistics summary location <active_RP>

Per PCB:

show tcp brief | i ":179"

show tcp detail pcb <pcb> location 0/RP0/CPU0

show tcp statistics pcb <pcb> location <active_RP>

Global TCP statistics commands show the overall health of TCP sessions. Apart from the data packet statistics (in/out), you can see for example whether there are packets with checksum errors, malformed packets, packets dropped due to authentication errors, out-of-order packets, packets with data after window, which gives you an indication of the behaviour of TCP peers.

In the per-PCB commands, you can see important parameters of a TCP session, like MSS, maximum roundtrip time, and so on.

Relevant counters in the output of show tcp detail pcb command are:

- **Retrans Timer Starts**: indicates how many times the retransmission timer started.
- Retrans Timer Wakeups: indicates how many times did the retransmission timer run out, triggering a retransmission of the TCP segment.
- **Current send queue size in bytes**: unacknowledged bytes from the peer.
- **Current receive queue size in bytes/packets**: bytes/packets yet to be read by the application (BGP).
- **mis-ordered bytes**: bytes that are queued in the save queue due to a hole in TCP receive window.

 $<\#root$

RP/0/RSP0/CPU0:ASR9k-B#

show tcp detail pcb 0x4a4400e4

==

Connection state is ESTAB, I/O status: 0, socket status: 0 Established at Sat Apr 20 18:26:26 2024

PCB 0x4a4400e4, SO 0x4a42c0ac, TCPCB 0x4a43b708, vrfid 0x60000000, Pak Prio: Normal, TOS: 64, TTL: 255, Hash index: 402 Local host: 10.10.10.229, Local port: 179 (Local App PID: 856311) Foreign host: 10.10.10.254, Foreign port: 46980 (Local App PID/instance/SPL_APP_ID: 856311/0/0)

Current send queue size in bytes: 0 (max 16384)

Current receive queue size in bytes: 0 (max 65535)

mis-ordered: 0 bytes

Current receive queue size in packets: 0 (max 60)


```
SRTT: 162 ms, RTTO: 415 ms, RTV: 253 ms, KRTT: 0 ms
minRTT: 0 ms, maxRTT: 247 ms
ACK hold time: 200 ms, Keepalive time: 300 sec, SYN waittime: 30 sec
Giveup time: 0 ms, Retransmission retries: 0, Retransmit forever: FALSE
Connect retries remaining: 0, connect retry interval: 0 secs
\langle \ldots \rangleRP/0/RSP0/CPU0:ASR9k-B#
```
Monitor Memory Utilisation

The BGP route table is stored in BGP process heap memory. The routing table is stored in RIB process heap memory.

Useful commands for heap memory monitoring:

show memory summary

show memory summary detail

show memory-top-consumers

show memory heap summary all

Telemetry sensor path:

Cisco-IOS-XR-nto-misc-oper:memory-summary/nodes/node/detail

FIB stores forwarding entries in shared memory space.

Useful commands for shared memory monitoring:

 show memory summary show memory summary detail show shmwin summary

Monitor BGP Process Performance

Useful command that provides internal data on BGP process performance:

```
  show bgp process performance-statistics
```
Monitor BGP Convergence

Another useful command is the one that shows the overall status of BGP convergence: show bgp convergence

When the network is stable, you see something like this:

RP/0/RP0/CPU0:ASR9k-B#show bgp convergence Mon Dec 18 13:55:47.976 UTC Converged. All received routes in RIB, all neighbors updated. All neighbors have empty write queues. RP/0/RP0/CPU0:ASR9k-B#