BGP RR Scale Considerations and KPI Monitoring

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Introduction

This document describes the major contributors to the maximum scale a Border Gateway Protocol (BGP) Route-Reflectors (RR) can achieve and guidance on BGP RR performance monitoring.

HW/SW Platform Selection

A high-scale BGP RR is typically not in the forwarding path of packets carrying services provided by an Internet Service Provider. Therefore, the hardware requirements for a BGP RR and routers that are predominantly forwarding packets in the data path are different. Standard routers are built with a powerful data-path forwarding element and a comparatively moderate control-path element. A BGP RR performs all its tasks in a control plan.

Physical Cisco IOS XR Router	Cisco IOS XRv 9000 Appliance	Cisco IOS XRv 9000 Router (aka XRv9k)
 Moderate control plane capacity (typically between 2 and 6 CPU cores allocated to RP XR VM) Unused data-path capacity 	 High control plane capacity (on Cisco UCS M5-based appliance 36 CPU cores are dedicated to RP XR VM) Equal split between data- path and control-path capacity. XRv9k image runs on barebone for maximum performance 	 Customisable control plane capacity Equal split between data- path and control-path power when using BGP RR image. An additional layer of virtualisation impacts performance.

Within the Cisco IOS® XR family of products, you can choose between 3 flavours of HW/SW platforms for a BGP RR role:

As of this writing, XRv9k Appliance is the optimal platform choice for BGP RR because it provides for highest control plane capacity with maximum performance.

Scale And Performance Considerations

The supported scale of data-plane entities is relatively easy to express because the performance of the datapath element seldom depends on the scale. For example, a TCAM lookup takes the same time regardless of the number of active TCAM entries.

The supported scale of control-plane entities is often much more complex because the scale and performance are interconnected. Consider a BGP RR with 1M routes. The work that a BGP process must carry out to maintain this BGP table depends on:

- 1. How many BGP peers are active?
- 2. What address families are active?
- 3. How are they distributed into update groups?
- 4. The complexity of RPLs (Route Policies)
- 5. Frequency of updates (incoming updates and also outgoing updates advertisement interval).
- 6. TCP MSS, Interface/Path MTU tuning this will aid in better performance
- 7. If dual-RP, is NSR enabled
- 8. Any known slow-peers, that aren't in separate update-group
- 9. Nexthop trigger-delay value

Number Of BGP Peers

The number of BGP peers is usually the first and unfortunately, often the only thing that comes to mind when considering the BGP scale. While the supported BGP scale cannot be represented without mentioning the number of BGP peers, it is not the most important factor. Many other aspects are equally relevant.

Address Families

The type of address family (AF) is an important factor in BGP performance considerations because in typical deployments it impacts the size of a single route. The number of IPv4 routes that can be packed into a single TCP segment is significantly higher than the number of VPNv4 routes. Hence for the same scale of BGP table changes, an IPv4 BGP RR has less work to do compared to a VPNv4 BGP RR. Obviously, in deployments where a significant number of communities is added to every route, the difference between AFs becomes less significant, but the size of a single route is then even bigger and requires consideration.

Number Of Update Groups

The BGP process prepares a single update for all members of the same update group. Then TCP process splits the update data into a required number of TCP segments (depending on TCP MSS) towards each member of the update group. You can see the active update groups and their members by using the show bgp update-group command. You can influence which and how many peers are members of an update group by creating a common outbound policy for a group of peers you want to be in the same update group. A single update sent by the BGP RR to a high number of BGP RR clients can trigger a burst of TCP ACKs that can be dropped in Local Packet Transport Service (LPTS) component of Cisco IOS XR routers.

Complexity of RPLs (Route Policies)

The complexity of route policies used by BGP impacts the BGP process performance. Every received or sent route must be evaluated against the configured route policy. A very long policy requires many CPU cycles to be spent on this action. A route policy that includes a regular expression is especially heavy on processing. A regular expression helps you express the route policy in a lesser number of lines but requires more CPU cycles while processing than the equivalent route policy that does not use regular expression.

Frequency Of Updates

The frequency of updates has an important bearing on the BGP scale. The number of updates is often difficult to predict. You can influence the frequency of updates by using the "advertisementinterval" command, which sets the minimum interval between the sending of BGP) routing updates. The default value towards iBGP peers is 0 seconds and 30 towards eBGP peers it is 30 seconds.

TCP MSS And Interface/Path MTU

Splitting an update into many TCP segments can put a lot of strain on TCP process resources in a high-scale and high-update frequency environment. A bigger path MTU and bigger TCP MSS are better for BGP and TCP performance.

NSR On Dual-RP Routers

NSR is a great feature for redundancy, but it does have an impact on BGP performance. On Cisco IOS XR routers both RPs are receiving simultaneously every BGP update directly from the NPU on the ingress line card, which means that the Active RP does not have to spend time on replicating the update to the Standby RP. However, every update generated by the Active RP must be sent to the Standby RP and from there to the BGP peer. This allows the Standby RP to be always up to date on the sequence and acknowledgement numbers but does have an impact on the overall BGP performance. This is why it is recommended that a BGP RR is a single-RP router.

Slow Peers

A slow peer can slow down the updates towards all members of the update group because the BGP process must keep the update in its memory until all peers have acknowledged it. If you know that some peers are much slower (for example routers in a legacy part of the network), separate them upfront into an update group. By default, Cisco IOS XR reports a slow peer via syslog message. You can create static slow peers (that never share the update group with others) or fine-tune the dynamic slow peer behaviour by using the BGP slow-peer configuration command in global or per-neighbour config mode. A good further read on this can be found in <u>Troubleshoot Slow BGP Convergence Due to Suboptimal Route Policies on IOS-XR</u> on the Cisco xrdocs.io portal.

Nexthop trigger-delay

If multiple BGP next-hops change in a short time interval and the critical nexthop trigger-delay value of zero is configured in an address family (AF) with a high number of routes, a full walk of the AF must be executed on every next-hop change event. Repeated walks of that AF increase the convergence time in address families with lower critical nexthop trigger-delay values. You can see the next-hop trigger-delay values by running the "show bgp all all nexthops" command.

Example Of Validated Multidimensional BGP RR Scale

Multidimensional scale results, especially for control plane features, are highly dependent on the specific test environment. Test results can vary significantly if some of the parameters are changed.

Parameter	Value	Value
Platform	XRv9k Appliance (UCS M5 based)	ASR9902
IOS XR release	7.5.2 + umbrella SMU for Cisco bug ID <u>CSCwf09600</u> (constituents of this umbrella SMU are integrated in Cisco IOS XR release 7.9.2 and later)	7.11.2
Peers	VPNv4 eBGP: 2500 VPNv4 iBGP: 1700	VPNv4 iBGP: 2000
BGP Routes	Per session: 200 Total: 400k	Per session: 750 VPNv4: 1.36M

IGP Routes	Paths per route: 1 10k (ISIS)	VPNv6: 150k IPv4: 950k IPv6: 200k Total: ~2.6M Paths per route: 1 10k (ISIS)
BGP update groups	1	1
BGP Timers	default	default
LPTS BGP-known policer rate	50,000	25,000
tcp num-thread configuration	16 16	16 16
BGP send-buffer-size	default	default
Key Performance Indicators (KPI) summary	 Test case with highest input and output packet rate: Input: 49.4kpps Output: 95kpps ==> LPTS drops (policer at 50kpps) ==> No drops in NetIO clients ==> Max XIPC Queue size (BGP): 1362 ==> Max XIPC Queue size (TCP): 1248 	 Test case with highest input packet rate: Input: 16030 pkts/s Output: 31 pkts/s ==> No drops in LPTS nor NetIO clients ==> Max XIPC Queue size (BGP): 378 ==> Max XIPC Queue size (TCP): 1021 Test case with highest output packet rate: Input: 12172 pkts/s Output: 23465 pkts/s ==> No drops

		in LPTS nor NetIO clients ==> Max XIPC Queue size (BGP): 109 ==> Max XIPC Queue size (TCP): 1518
--	--	---

Design Considerations

There are two approaches to BGP RR placement in the network:

- 1. Centralised/Flat BGP RR design.
- 2. Distributed/Hierarchical BGP RR design.

In a centralised/flat design, all BGP RR clients in the network establish BGP peering with a set (usually a pair) of BGP RR devices that hold the exact same information. This approach is simple to implement and works well in small to moderate scale networks. Any change in BGP table is propagated quickly to all BGP RR clients. As the number of BGP RR clients grows, the design can reach a scale limit when the number of TCP connections on the BGP RR devices grows to the extent where their performance is impacted.

In a distributed/hierarchical design, network is split into several regions. All routers in a region establish BGP peering with a set (usually a pair) of BGP RR devices that hold the exact same information. These BGP RR devices act as BGP RR clients to another set (usually a pair) of BGP RR devices. This design approach allows for easy network expansion, while keeping the number of TCP connections on every single BGP RR under certain limit.

Another design consideration is tailoring the scope of recipients of BGP updates. Depending on the VRF distribution among BGP RR clients, it's worth considering the RT Constrained Route Distribution. If all BGP RR clients have interfaces in the same VRF, RT Constrained Route Distribution does not bring many benefits. However, if VRFs are sparsely distributed among all BGP RR clients, use of RT Constrained Route Distribution significantly reduces *f* the load on the bgp process on the BGP RR.

Monitor BGP Key Performance Indicators (KPI)

Monitoring of BGP RR's Key Performance Indicators (KPI) is important for ensuring proper network operation.

A significant change in the network topology (for example a major DWDM link flap) can trigger routing updates that generate excessive traffic towards and/or from the BGP RR. Significant traffic hitting the BGP RR typically carries:

- 1. Updates from BGP peers.
- 2. TCP ACKs generated by the BGP peers, in response to updates sent by BGP RR & vice-versa

This section of the document explains the KPI that need to be monitored on a typical BGP RR and also how to tell which of the two significant BGP traffic types are causing high control plane traffic rate.

Path of BGP packets inside the router can be depicted as follows:

Punt

Ethernet controller -(packet)-> datapath forwarder -(packet)-> LPTS -(packet)-> SPP -(packet) -> NetIO -(packet)-> TCP -(message)-> BGP

Inject

BGP -(message)-> TCP -(packet)-> NetIO -(packet)-> SPP -(packet) -> datapath forwarder -(packet)-> Ethernet controller

KPIs can be split into:

Essentials:

- 1. Datapath Forwarder
- 2. LPTS (hardware punt policers settings, accept counters and drop counters)
- 3. SPP
- 4. NetIO
- 5. IPC queues (NetIO <==> TCP <==> BGP)
- 6. BGP InQ/OutQ sizes

Optional:

- 1. CPU utilisation
- 2. Memory utilisation
- 3. TCP statistics
- 4. BGP process performance
- 5. BGP convergence

Monitor Datapath Forwarder

On XRv9000 the datapath forwarder is the Data Plane Agent (DPA), while on ASR9000 platforms it is the Network Processor (NP).

Monitor XRv9000 Data Plane Agent (DPA)

Useful command to see the load and statistics of the DPA is:

```
show controllers dpa statistics global
```

This command shows all the non-zero counter, that give you insight into the type and number of packets punted from network interfaces to RP CPU, injected from RP CPU towards network interfaces, and the numer of packets dropped:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show controllers dpa statistics global

Index	Debug	Count
350	TBPG L2 mailbox events	1
Index	Punt	Count
1500 1503 1611 1776 2001	CDP ARP IFIB LLDP IPv4 incomplete Tx adjacency	46790 212 595305 94037 2
Index	Inject	Count
744 749 765 766	CLNS multicast from fabric pre-route IPv4 from fabric Inject to fabric Inject to port	321250 273993 595245 141033
Index	Drop	Count
416 474 673	Egress UIDB in down state IPv4 egress UIDB down Pre-route PIT lookup missed	1 2 2

Monitor ASR9000 Network Processor (NP)

Useful commands to see the load and statistics of each NP in the system are:

show controllers np load all show controllers np counters all

NP on ASR9000 has a rich set of counters that show you the number, rate and type of processed and dropped packets,.

<#root>

RP/0/RSP0/CPU0:ASR9k-B#

show controllers np load all

Node: 0/0/CPU0:

Load Packet Rate NPO: 0% utilization 937 pps NP1: 0% utilization 538 pps RP/0/RSP0/CPU0:ASR9k-B# RP/0/RSP0/CPU0:ASR9k-B#

show controllers np counters all

Node: 0/0/CPU0:

Show Last Read	global stats counters for NPO, revision v4 clearing of counters for this NP: NEVER 92 non-zero NP counters:		
Offse	et Counter	FrameValue	Rate (pps)
			10
10	MDF_IX_LC_CPU	25475508	10
17	MDF_TX_WIRE	681957877	267
21	MDF_TX_FABRIC	683500690	267
33	PARSE_FAB_RECEIVE_CNT	681767730	267
37	PARSE_INTR_RECEIVE_CNT	1323024637	517
41	PARSE_INJ_RECEIVE_CNT	13949634	5
45	PARSE_ENET_RECEIVE_CNT	677655725	265
49	PARSE_TM_LOOP_RECEIVE_CNT	49331192	19
53	PARSE_TOP_LOOP_RECEIVE_CNT	1520846	1
109	RSV_DROP_EGR_UIDB_NO_MATCH	10	0
146	RSV_DROP_IPV4_RXADJ_DROP	1	0
164	RSV_DROP_ING_LAG_NO_MATCH	3	0
241	RSV_DROP_MPLS_LEAF_NO_MATCH	1312	0
<	.>		

```
Monitor LPTS
```

As a standard BGP RR is not in the forwarding path, all packets received on network interface are punted to control-plane. The data-path element on a BGP RR performs a small number of simple operations before packets are punted to control-plane. As the data-path element is unlikely to be a point of congestion, the only element on the line card that needs monitoring are the LPTS stats.

Please note that in case of XRv9k, hardware statisitics map to the vPP

Command:

show lpts pifib hardware police location <location> | inc "Node|flow_type|BGP"

Example:

RP/0/RP0/CPU0	:xrv9k-01	#sh lpts pifik	o hardware police	e location ()/0/CPI	JO i "Node flo	ow_type BGP"	
	Node 0/0/0	CPUO:						
flow_type	priority	<pre>sw_police_id</pre>	hw_policer_addr	Cur. Rate	burst	<pre>static_avgrate</pre>	avgrate_type	AggrAc
BGP-known	high	6	220	50000	1250	2500	Global	164013
BGP-cfg-peer	medium	7	221	4000	1000	2000	Global	355976
BGP-default	low	8	222	3000	750	1500	Global	521238
RP/0/RP0/CPU0	:xrv9k-01	#						

What to look for: If a significant jump in AggDrops against BGP-known flow type is observed, start looking for network topology changes that have triggered such massive control plane churn.

Telemetry data path:

Cisco-IOS-XR-lpts-pre-ifib-oper:lpts-pifib



Note: LPTS stats counters can be cleared. Your monitoring system must account for this possibility.

Monitor SPP

SPP is the first entity on the route processor or line card CPU that receives the packet punted from the NP or DPA via internal fabric, and the last point in the software packet processing before it is handed over to the fabric for injection into the NP or DPA.

Relevant commands for SPP monitoring:

show spp node-counters

show spp client

The show spp node-counters command shows the rate of punted/injected packets and is easy to read and understand. For BGP sessions the relevant coounters are under client/punt and client/inject on the active RP.

6

92

The show spp client is more rich in output and gives a more detailed insight into number of packets enqueued/dropped towards clients, as well as the high watermark.

<#root> RP/0/RP0/CPU0:xrv9k-01# show spp node-counters 0/RP0/CPU0: socket/rx Punted packets: 595305 Punt bulk reads: Punt non-bulk reads: 595293 Management packets: 74200158 Management bulk reads: 1775930 Management non-bulk reads: 70031734 ----socket/tx Injected packets: 595245 Management packets: 139939168 ----xrv9k/classify Forwarded to SPP clients: 74795463 -----client/inject Injected from client: 140534413 Non-bulk injects: 140534413 ----client/punt punted to client: 74795371 no client found - send to defa: 0/0/CPU0: <. . .>

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show spp client

Sat Apr 20 17:11:40.725 UTC 0/RP0/CPU0: Clients _____

<. . .> netio, JID 254 (pid 4591) Reconnect Pending: F, Exited: F, Keep Queues: F, Pakman Client: T Quota Current: 0, Limit: 16384, Drops 0 Queues: Key Cur Max Enqueues High Watermark (Timestamp) Drops
 0xffffffff
 0
 10
 0

 0x03000006
 0
 2048
 0
 0 (22:13:52.195 Feb 21 24 UTC) 0 0 0 (22:13:52.196 Feb 21 24 UTC) 0 0x030000060 204800 (22:13:52.196 Feb 21 24 UTC)0x030000070 30724148811 (23:03:30.721 Feb 21 24 UTC)0x030000080 102451 (13:41:28.389 Mar 13 24 UTC)0x030000090 20481804111 (23:03:31.565 Feb 21 24 UTC) 0 0 0

Monitor NetIO

While the LPTS policer only shows the count of packets accepted or dropped by a corresponding policer, at NetIO level we can see the rate of packets punted to RP CPU. Since on a typical BGP RR the great majority of received packets are BGP packets, the overall NetIO rate indicates very closely the rate of received BGP packets.

<#root>

Command:

show netio rates

Example:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show netio rates

Netio packet rate for node 0/RP0/CPU0

Current rate	(updated 0 seconds ago):	
Input:	7845 pkts/s	
Output:	10570 pkts/s	
Driver Output:	10598 pkts/s	
1 minute rate	(updated 0 seconds ago):	
Input:	7825 pkts/s	
Output:	10542 pkts/s	
Driver Output:	10569 pkts/s	
5 minute rate	(updated 0 seconds ago):	
Input:	7997 pkts/s	
Output:	10677 pkts/s	
Driver Output:	10704 pkts/s	

RP/0/RP0/CPU0:xrv9k-01#

What to look for:

• If a significant jump in NetIO rate is observed, start looking for network topology changes that have triggered such massive control plane churn.

Telemetry data path:

• not applicable as telemetry must stream counter values, not rates. BGP-known LPTS policer accept counter can be used on the Telemetry collector to approximate the average rate of received BGP packets from known peers.

Monitor XIPC Queues

On the punt path packets received by NetIO from LPTS are passed on to TCP and BGP. It's important to monitor these queues:

- 1. TCP high priority queue though which NetIO delivers packets to TCP
- 2. BGP control queue
- 3. BGP data queue

On the inject path packets are created by TCP and passed onto NetIO. It's important to monitor these queues:

1. OutputL XIPC queue

Commands:

show netio clients show processes bgp | i "Job Id" show xipcq jid <bgp_job_id>

show xipcq jid
dp_job_id> queue-id <n>

Examples:

NetIO to TCP, view from NetIO standpoint:

RP/0/RP0/CPU0:xrv9k-01#show netio clients

NE 1 1 2				
	Input	Punt	XIPC InputQ	XIPC PuntQ
ClientID	Drop/Total	Drop/Total	Cur/High/Max	Cur/High/Max
<>				
tcp	L 0/340774 H 0/44774091	0/0	L 0/10/12800 H 0/784/12800	0/0/0

TCP to NetIO, view from NetIO standpoint ::

XIPC queues	Dropped/Queued	Cur/High/Max
OutputL	124860/9355257	0/14000/14000

NetIO to TCP, view from TCP process standpoint:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show processes tcp

| i "Job Id"

Job Id: 430 RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

430

Mon Apr 17 16:16:11.315 CEST τд Namo

Id	Name	Size	Cur Size	e Produced	Dropped	HWM
17	XIPC_xipcq_12_0_9854_6506_i	60000	0	39010269	0	960
16	XIPC_xipcq_12_0_9854_6506_i	24000	0	31518436	0	1364
15	XIPC_tcp_124	3200	0	0	0	0
14	XIPC_tcp_125	9600	0	0	0	0
13	XIPC_tcp_psb_0	25600	0	0	0	0
10	XIPC_tcp_iq_9	102400	0	9486010	0	312
12	XIPC_tcp_iq_8	102400	0	8892274	0	280
9	XIPC_tcp_iq_5	102400	0	8291392	0	289
11	XIPC_tcp_iq_7	102400	0	9700123	0	319
8	XIPC_tcp_iq_6	102400	0	9378703	0	332
7	XIPC_tcp_iq_3	102400	0	7221706	0	261
6	XIPC_tcp_iq_4	102400	0	9791070	0	308
4	XIPC_tcp_ctrl_iq_1	4266	0	0	0	0
5	XIPC_tcp_iq_2	102400	0	9699027	0	295
3	XIPC_tcp_ctrl_iq_0	4266	0	209909	0	9
2	XIPC_tcp_i1	12800	0	39460564	0	784
1	XIPC_tcp_i0	12800	0	212540	0	10

TCP to BGP:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show processes bgp

| i "Job Id"

Job Id: 1078 RP/0/RP0/CPU0:xrv9k-01#

Name

show xipcq jid

1078 Mon Apr 17 16:09:33.046 CEST

Id

Dropped

HWM

2 XIPC_xipcq_12_0_9854_6506_i	60000	2	35546667	0	15034
1 XIPC_xipcq_12_0_9854_6506_i	24000	0	1369029	0	50
D / O / D D O / C D U O D W m O U O 1 H					

RP/0/RP0/CPU0:xrv9k-01#

```
BGP data queue:
```

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

1078

queue-id 1

XIPC_xipcq_12_0_9854_6506_inst_1_data_toapp

```
:
```

Magic:	12344321
Version:	0
SHM Size:	192392
Owner PID:	9854
Owner JID:	1078
Queue ID:	1
Owner MQ handle:	483
User Context:	0x64
Interrupt Flag:	0
Sent-on-full Flag:	0
Max Queue Size:	24000
Queue Size:	0
Client Queue Size:	24000
High watermark:	50
Last Trigger Sent:	Mon Apr 17 16:11:10.009 CEST
MQ Send Errors:	0
Priority Queues:	
Duri e Cita	During Tratal

Prio	Size	Drops	lotal			
Unspec	24000	0	0			
Normal	24000	0	1396159			
Medium	24000	0	0			
High	24000	0	0			
Crucial	24000	0	0			
RP/0/RP0/CPU0:xrv9k-01#						

BGP control queue:

<#root>

RP/0/RP0/CPU0:xrv9k-01#

show xipcq jid

1078

queue-id

2

XIPC_xipcq_12_0_9854	↓_6506_ir	nst_1_c	trl_toapp:				
Magic:	12344322	1					
Version:	0						
SHM Size:	480392						
Owner PID:	9854						
Owner JID:	1078						
Queue ID:	2						
Owner MQ handle:	485						
User Context:	0x64						
Interrupt Flag:	0						
Sent-on-full Flag:	0						
Max Queue Size:	60000						
Queue Size:	0						
Client Queue Size:	60000						
High watermark:	15034						
Last Trigger Sent:	Mon Apr	17 16:	12:49.483	CEST			
MQ Send Errors:	0						
Priority Queues:							
Prio Siz	ze [Drops	Total				
Unspec 6000)0	0	0				
Normal 6000	00	Õ	37313633				
Medium 6000	0	Õ	0				
High 6000)0	õ	0				
Crucial 6000)0	õ	0				
RP/0/RP0/CPU0:xrv9k-	-01#	Ŭ	Ũ				
	• = "						

What to look for:

- there must be no drops in relevant queues
- in XIPC queue stats High watermark (HWM) must not exceed the 50% of queue size

For better tracking of high watermark value evolution, you must clear the high watermark value after every read. Note that this does not clear only the HWM counter but also clears all queue statistics. The format of the command for clearing the XIPC queue statistics is: clear xipcq statistics queue-name <queue_name>

As the queue name often includes the Process ID (PID), the queue name changes after process restart. Some examples of commands for clearing the relevant queues statistics:

```
clear xipcq statistics queue-name XIPC_tcp_i0
clear xipcq statistics queue-name XIPC_tcp_i1
clear xipcq statistics queue-name XIPC_xipcq_12_0_9854_6506_inst_1_data_toapp
clear xipcq statistics queue-name XIPC_xipcq_12_0_9854_6506_inst_1_ctrl_toapp
```

Telemetry path:

• There are no Telemetry sensor paths for XIPC.

Monitor BGP Input And Output Queues

BGP maintains an input and output queue for every BGP peer. Data sits in InQ when TCP has passed it to BGP, but BGP hasn't processed them yet. Data sits in OutQ while BGP waits on TCP to split the data into packets and transmit them. The instantaneous size of BGP InQ/OutQ provides a good indication of how busy the BGP process is.

Command:

```
show bgp <AFI> <SAFI> summary
```

Example:

RP/0/RP0/CPU0:xrv9k-01#show bgp all all summary
Address Family: VPNv4 Unicast

BGP router identifier 192.168.0.1, local AS number 65000 BGP generic scan interval 60 secs BGP table state: Active Table ID: 0x0 BGP main routing table version 2208096 BGP scan interval 60 secs

BGP is operating in STANDALONE mode.

Process	RcvTb	lVer	bRIB/RI	B Label	lVer	Impo	rtVer	Ser	ndTblVer	StandbyVer
Speaker	2208	8096	2208090	5 2208	3096	22	08096		2208096	2208096
Neighbor	Spk	AS	MsgRc∨d	MsgSent	Tbl	Ver	InQ	OutQ	Up/Down	St/PfxRcd
10.0.0.2	0	65000	180	601022	2208	3096	0	0	02:56:18	100
10.0.0.3	0	65000	180	601022	2208	3096	0	0	02:56:18	100
10.0.0.4	0	65000	180	601022	2208	3096	0	0	02:56:21	100
10.0.0.5	0	65000	180	601022	2208	3096	0	0	02:56:21	100
10.0.0.6	0	65000	180	601022	2208	3096	0	0	02:56:18	100

What to look for:

- The size of InQ/OutQ must be zero when the network is stable. It changes quickly when updates are exchanged.
- InQ/OutQ size must not monotonously increase over time.

Telemetry path:

• Cisco-IOS-XR-ipv4-bgp-oper:bgp

Monitor BGP Message Rates

Some BGP neighbors can continuously send updates or withdrawals if network topology is unstable. The BGP RR must then replicate such routing table change thousands of times to all its RR clients. Therefore it is important to monitor the message rates received from neighbors, to track sources of instabilities.

Command:

show bgp <AFI> <SAFI> summary

Example:

RP/0/RP0/CPU0:xrv9k-01#show bgp all all summary
Address Family: VPNv4 Unicast

BGP router identifier 192.168.0.1, local AS number 65000 BGP generic scan interval 60 secs BGP table state: Active Table ID: 0x0 BGP main routing table version 2208096 BGP scan interval 60 secs

BGP is operating in STANDALONE mode.

Process	RcvTb ⁻	lVer	bRIB/RI	3 Label	Ver 1	Impo	rtVer	' Ser	ndTblVer	StandbyVer
Speaker	2208	8096	2208096	5 2208	8096	22	08096	j	2208096	2208096
Neighbor	Spk	AS	MsgRc∨d	MsgSent	Tb۱۱	/er	InQ	OutQ	Up/Down	St/PfxRcd
10.0.0.2	0	65000	180	601022	22080	096	0	0	02:56:18	100
10.0.0.3	0	65000	180	601022	22080	096	0	0	02:56:18	100
10.0.0.4	0	65000	180	601022	22080	096	0	0	02:56:21	100
10.0.0.5	0	65000	180	601022	22080	096	0	0	02:56:21	100
10.0.0.6	0	65000	180	601022	22080	096	0	0	02:56:18	100

RR clients queues have roughly the same amount of MsgSent but some neighbors can have a number of MsgRcvd higher than others. You must capture multiple snapshots of this command to assess the rate of messages.

Once you have identified the offending peers, you can then go through other commands like show bgp neighbor <neighbor> detail and show bgp neighbor <neighbor> performance-statistics or show bgp recent-prefixes to try to understand which prefixes are flapping and whether it is always the same ones or different ones.



Note: The MsgRcvd and MsgSent counters are per-neighbor but not per address-family. So when running a command like show bgp all all summary, you see the same counters per neighbor in the sections for the various address-families. They do not represent the number of messages received/sent from/to that neighbor for that address-family but across address-families.

Monitor CPU Utilisation

CPU utilisation must be monitored on every router, but on a router with a high number of CPU cores dedicated to control plane some readouts can be unintuitive. On an BGP RR with a high number of CPU cores dedicated to Routing Processor (RP), as in the case of XRv9k Appliance, active threads run on different CPU cores, while a number of CPU cores remains idle. As a consequence some CPU cores can be very busy, but the overall CPU utilisation calculated across all CPU cores remains moderate.

Therefore, for proper monitoring of CPU cores utilisation via CLI, use the show processes cpu thread command.

Monitor TCP Statistics

Cisco IOS® maintains detailed statistics on every TCP session. CLI command show tcp brief displays the list of all existing TCP sessions. In this summary output, for every TCP session you can see this information:

- **PCB:** unique TCP session identifier.
- VRF-ID: the ID of the VRF in which the session exists.
 - To see the corresponding VRF name, run this command:
 - show cef vrf all summary | utility egrep "^VRF:|Vrfid" | utility egrep -B1 <VRF-ID>
- **Recv-Q**: instantaneous size of the receive Q. Receive queue holds packets received from NetIO. The **tcp** process extracts the data from a packet and sends it to the corresponding application.
- Send-Q: instantaneous size of the send Q. Send queue holds data received from an application. The tcp process splits the data into TCP segments (dictated by negotiated maximum segment size TCP MSS), encapsulates every segment into a layer 3 header of corresponding address family (IPv4 or IPv6) and sends the packet to NetIO.
- Local Address: local IPv4 or IPv6 address associated with the TCP socket. TCP sessions in LISTEN state are typically bound to "any" IP address, which is represented as "0.0.0.0" or "::" in case of IPv4 or IPv6 respectively.
- Foreign Address: remote IPv4 or IPv6 address associated with the TCP socket. TCP sessions in LISTEN state are typically bound to "any" IP address, which is represented as "0.0.0.0" or "::" in case of IPv4 or IPv6 respectively.
- State: TCP session state. Possible TCP session states are: LISTEN, SYNSENT, SYNRCVD, ESTAB, LASTACK, CLOSING, CLOSEWAIT, FINWAIT1, FINWAIT2, TIMEWAIT, CLOSED.

As the well-known BGP port number is 179, you can limit the displayed TCP sessions to those that are associated with the BGP application.

Example:

RP/0/RSP0/CPU0:ASR9	9k-B#show to	cp brief	incl	ude "PCB :179 "		
PCB	VRF-ID	Recv-Q	Send-Q	Local Address	Foreign Address	State
0x00007ff7d403bde0	0x6000000	0	0	:::179	:::0	LISTEN
0x00007ff7d403b020	0x6000002	0	0	:::179	:::0	LISTEN
0x00007ff7d403d130	0x6000000	0	0	192.168.0.4:50144	192.168.0.5:179	ESTAB
0x00007ff7a4025650	0x6000000	0	0	0.0.0.0:179	0.0.0:0	LISTEN
0x00007ff7a4024a50	0x6000002	0	0	0.0.0.0:179	0.0.0:0	LISTEN

You can use the displayed PCB value to obtain the statistics for a particular TCP session. CLI commands that provide insight into TCP process statistics:

Global:

show tcp statistics clients location <active_RP>

show tcp statistics summary location <active_RP>

Per PCB:

show tcp brief | i ":179"

show tcp detail pcb <pcb> location 0/RP0/CPU0

show tcp statistics pcb <pcb> location <active_RP>

Global TCP statistics commands show the overall health of TCP sessions. Apart from the data packet statistics (in/out), you can see for example whether there are packets with checksum errors, malformed packets, packets dropped due to authentication errors, out-of-order packets, packets with data after window, which gives you an indication of the behaviour of TCP peers.

In the per-PCB commands, you can see important parameters of a TCP session, like MSS, maximum round-trip time, and so on.

Relevant counters in the output of show tcp detail pcb command are:

- Retrans Timer Starts: indicates how many times the retransmission timer started.
- **Retrans Timer Wakeups**: indicates how many times did the retransmission timer run out, triggering a retransmission of the TCP segment.
- Current send queue size in bytes: unacknowledged bytes from the peer.
- Current receive queue size in bytes/packets: bytes/packets yet to be read by the application (BGP).
- **mis-ordered bytes**: bytes that are queued in the save queue due to a hole in TCP receive window.

<#root>

RP/0/RSP0/CPU0:ASR9k-B#

show tcp detail pcb 0x4a4400e4

Connection state is ESTAB, I/O status: 0, socket status: 0 Established at Sat Apr 20 18:26:26 2024

PCB 0x4a4400e4, S0 0x4a42c0ac, TCPCB 0x4a43b708, vrfid 0x6000000, Pak Prio: Normal, TOS: 64, TTL: 255, Hash index: 402 Local host: 10.10.10.229, Local port: 179 (Local App PID: 856311) Foreign host: 10.10.10.254, Foreign port: 46980 (Local App PID/instance/SPL_APP_ID: 856311/0/0)

Current send queue size in bytes: 0 (max 16384)

Current receive queue size in bytes: 0 (max 65535)

mis-ordered: 0 bytes

Current receive queue size in packets: 0 (max 60)

Timer	Starts	Wakeups	Next(msec)
Retrans	2795	0	0
SendWnd	1341	0	0
TimeWait	0	0	0
AckHold	274	2	0
KeepAlive	333	1	299983
PmtuAger	0	0	0
GiveUp	0	0	0
Throttle	0	0	0
FirstSyn	0	0	0

iss: 2030796738 snduna: 2034498828 sndnxt: 2034498828 sndmax: 2034498828 sndwnd: 3291 sndcwnd: 4200 irs: 285455091 rcvnxt: 285455710 rcvwnd: 64917 rcvadv: 285520627

```
SRTT: 162 ms, RTTO: 415 ms, RTV: 253 ms, KRTT: 0 ms
minRTT: 0 ms, maxRTT: 247 ms
ACK hold time: 200 ms, Keepalive time: 300 sec, SYN waittime: 30 sec
Giveup time: 0 ms, Retransmission retries: 0, Retransmit forever: FALSE
Connect retries remaining: 0, connect retry interval: 0 secs
<...>
RP/0/RSP0/CPU0:ASR9k-B#
```

Monitor Memory Utilisation

The BGP route table is stored in BGP process heap memory. The routing table is stored in RIB process heap memory.

Useful commands for heap memory monitoring:

show memory summary

show memory summary detail

show memory-top-consumers

show memory heap summary all

Telemetry sensor path:

Cisco-IOS-XR-nto-misc-oper:memory-summary/nodes/node/detail

FIB stores forwarding entries in shared memory space.

Useful commands for shared memory monitoring:

show memory summary show memory summary detail show shmwin summary

Monitor BGP Process Performance

Useful command that provides internal data on BGP process performance:

```
show bgp process performance-statistics
```

Monitor BGP Convergence

Another useful command is the one that shows the overall status of BGP convergence: show bgp convergence

When the network is stable, you see something like this:

RP/0/RP0/CPU0:ASR9k-B#show bgp convergence Mon Dec 18 13:55:47.976 UTC Converged. All received routes in RIB, all neighbors updated. All neighbors have empty write queues. RP/0/RP0/CPU0:ASR9k-B#